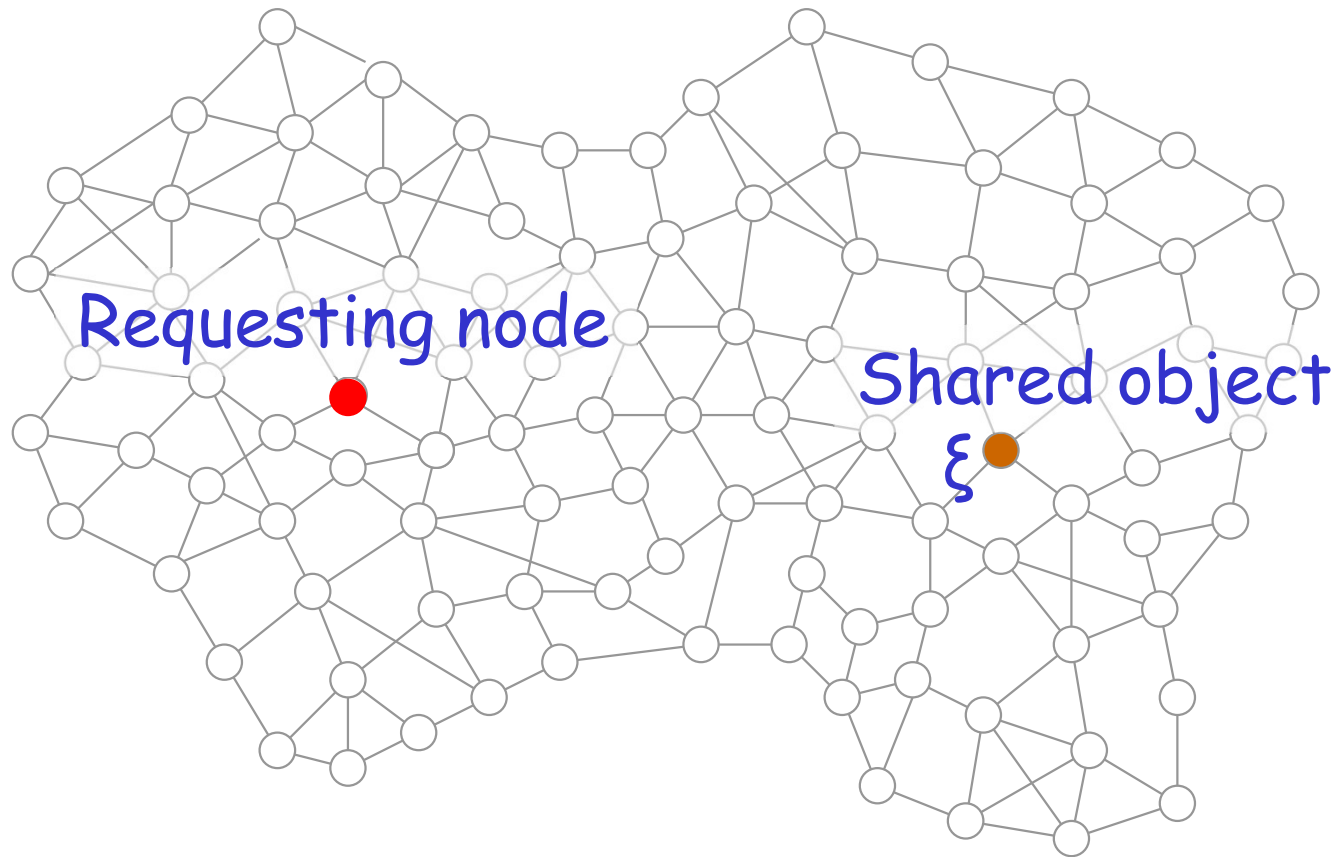


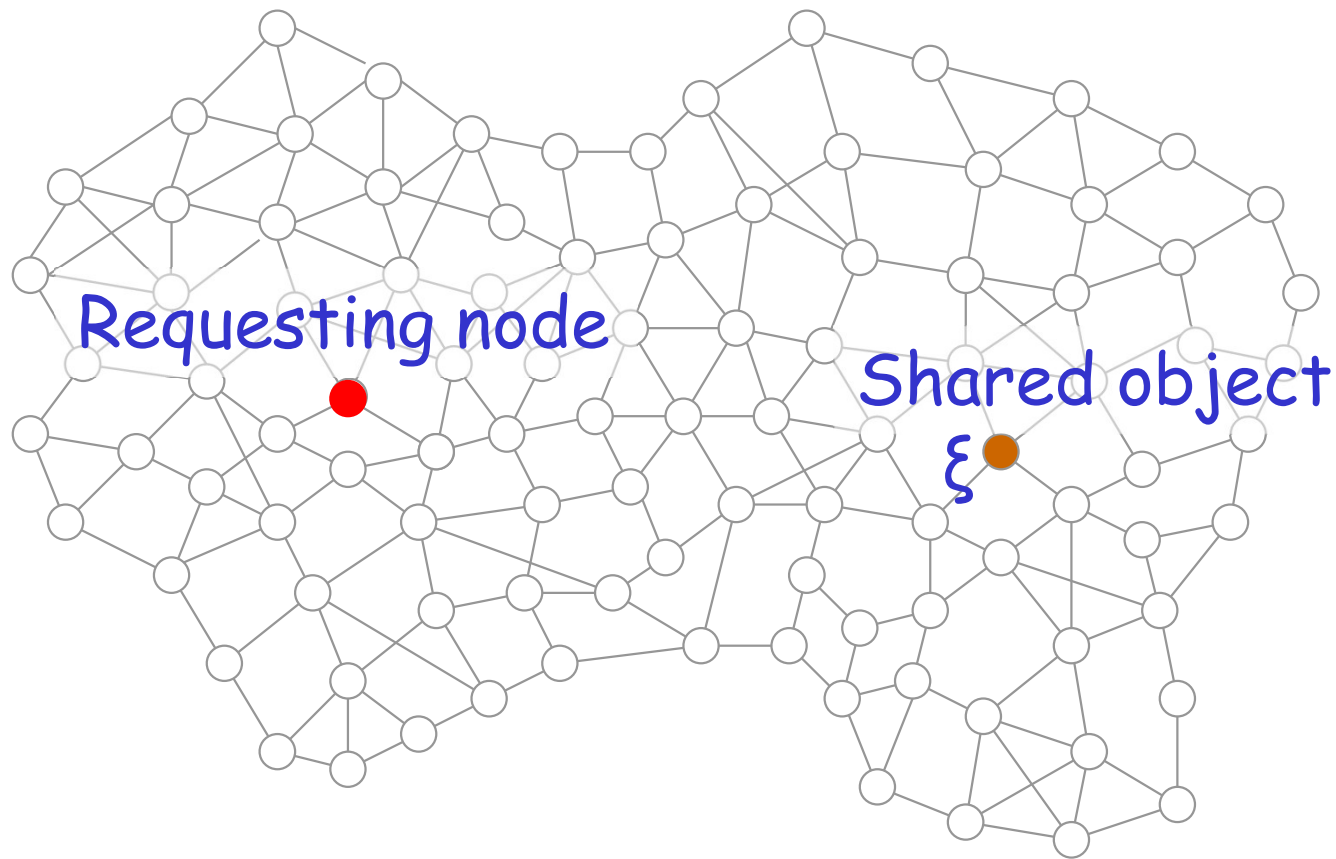
Learning-Augmented Distributed Directories

Swapnil Guragain, Bibek Maharjan, Sushant
Bhattarai, Gokarna Sharma, and Pavan Poudel

Kent State University, USA
University of Houston-Clear Lake, USA

November 5, 2025

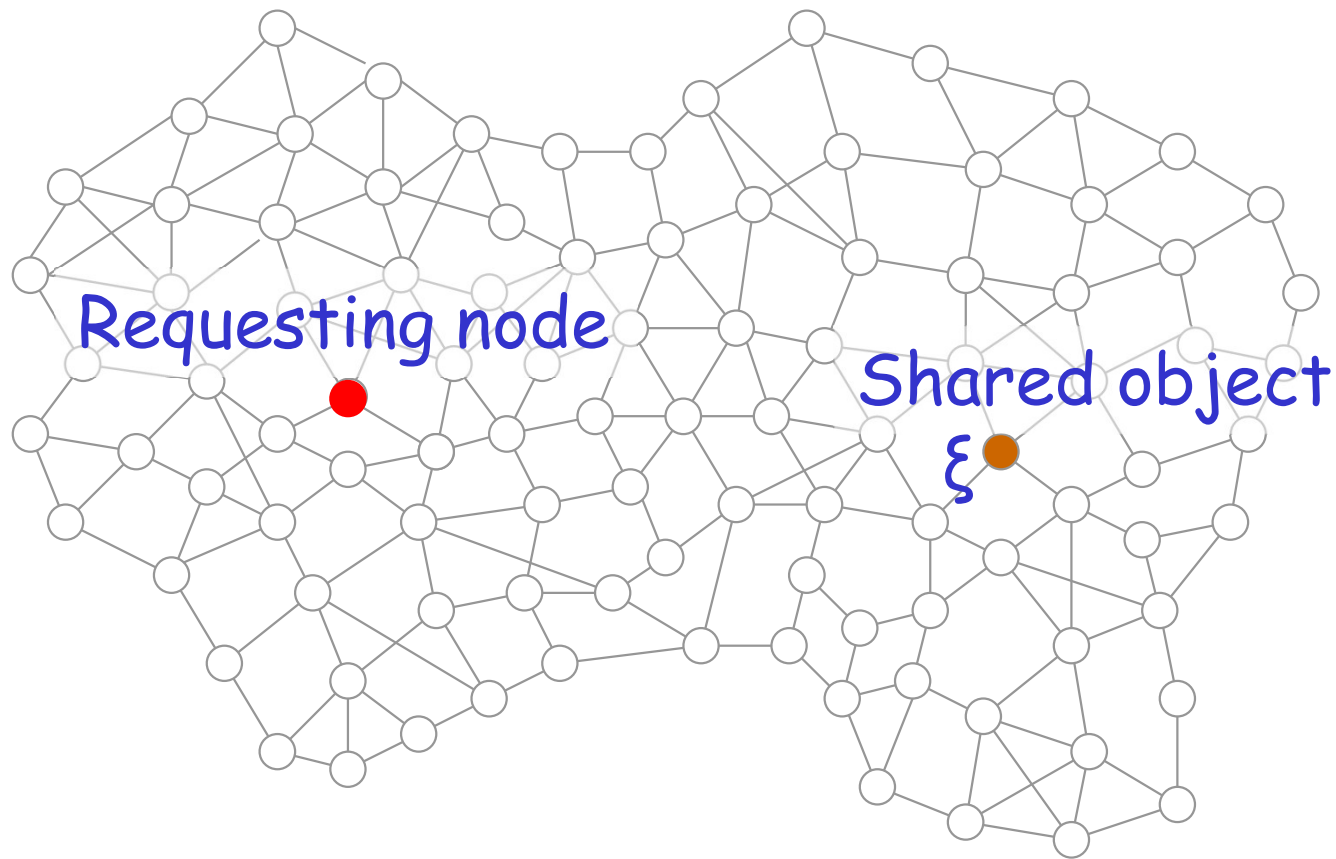




Naïve Implementation #1: Inform all other nodes

Update Cost: $O(m)$

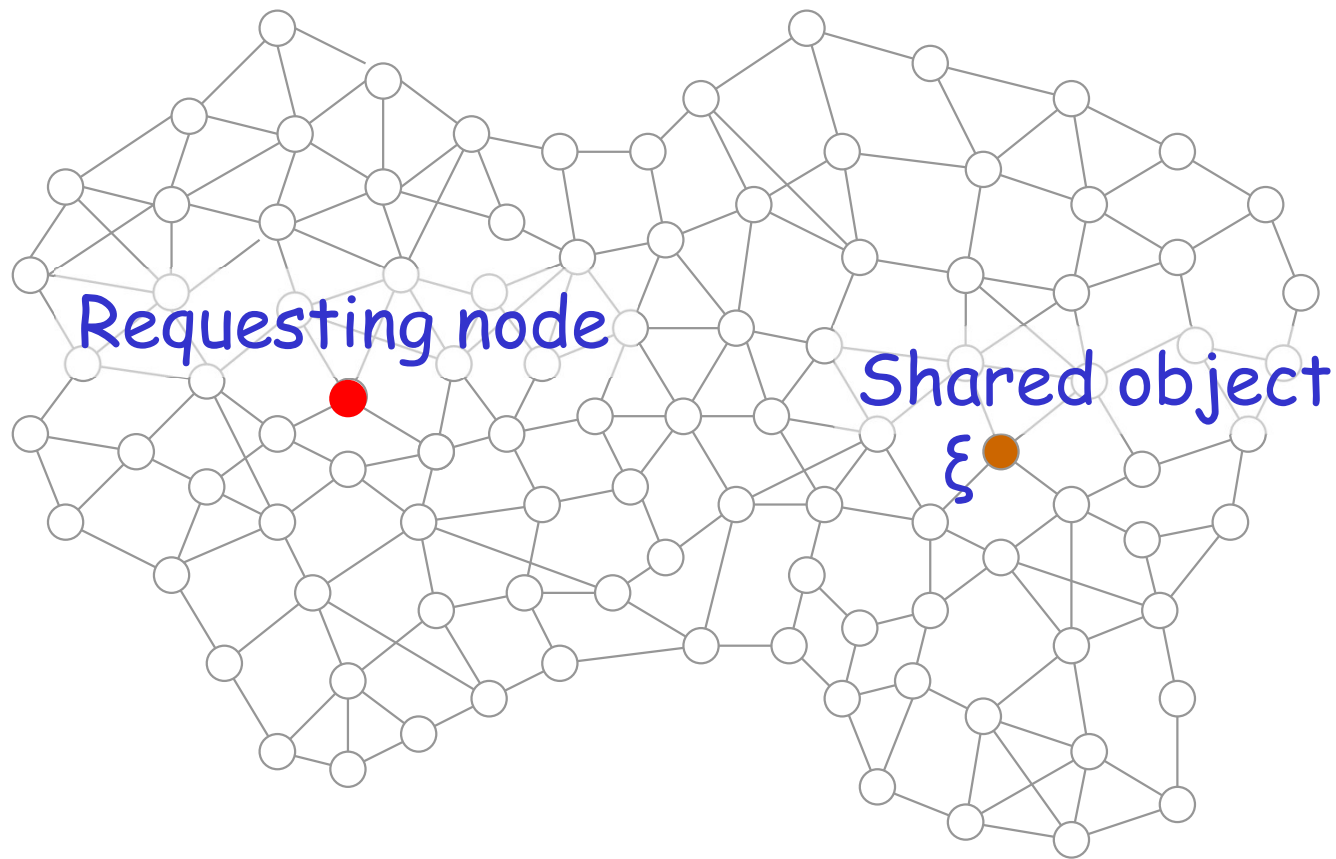
Lookup Cost: $O(1)$



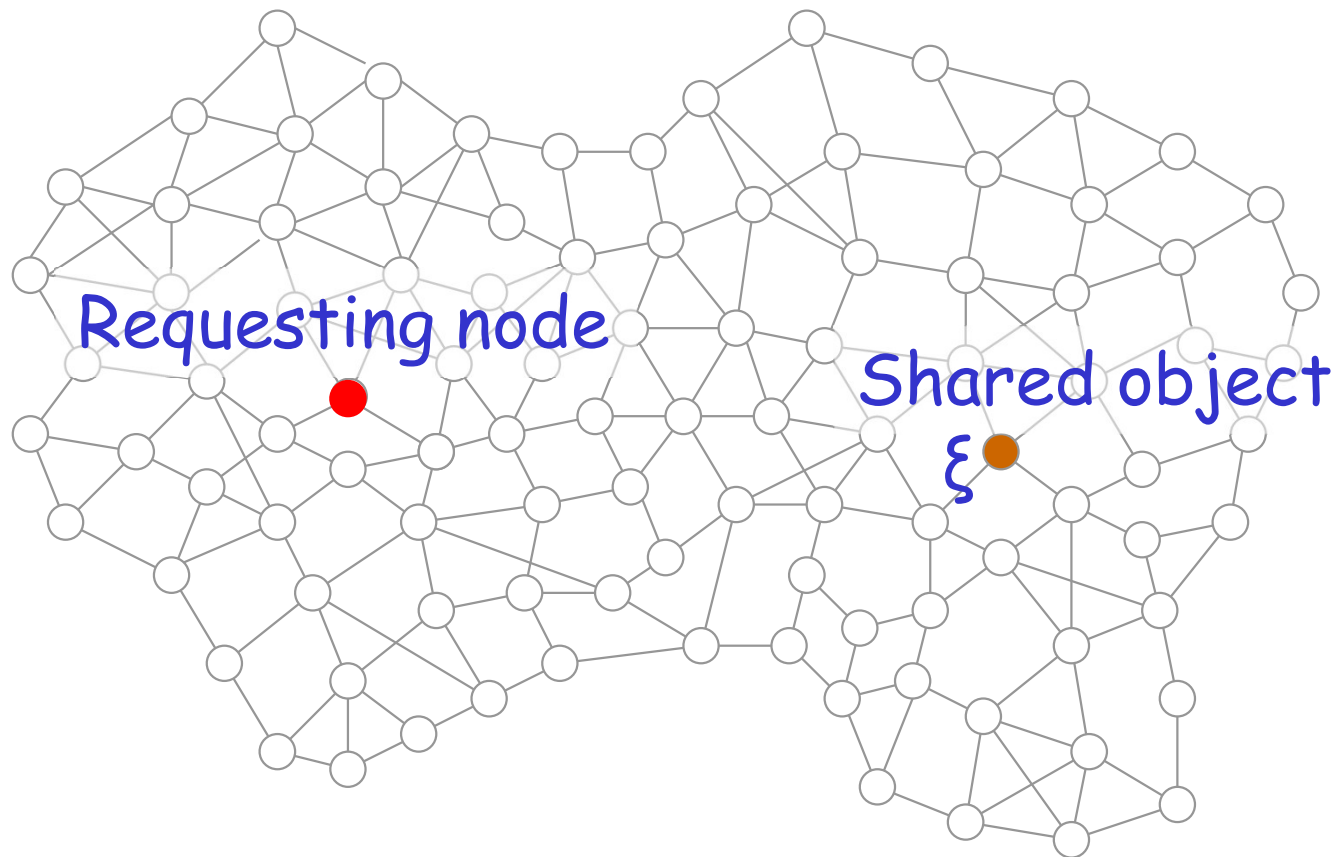
Naïve Implementation #2: Search the network

Update Cost: $O(1)$

Lookup Cost: $O(m)$



Question: Can both update and lookup costs be simultaneously minimized?



Optimal cost: Shortest path distance between requesting node and object node

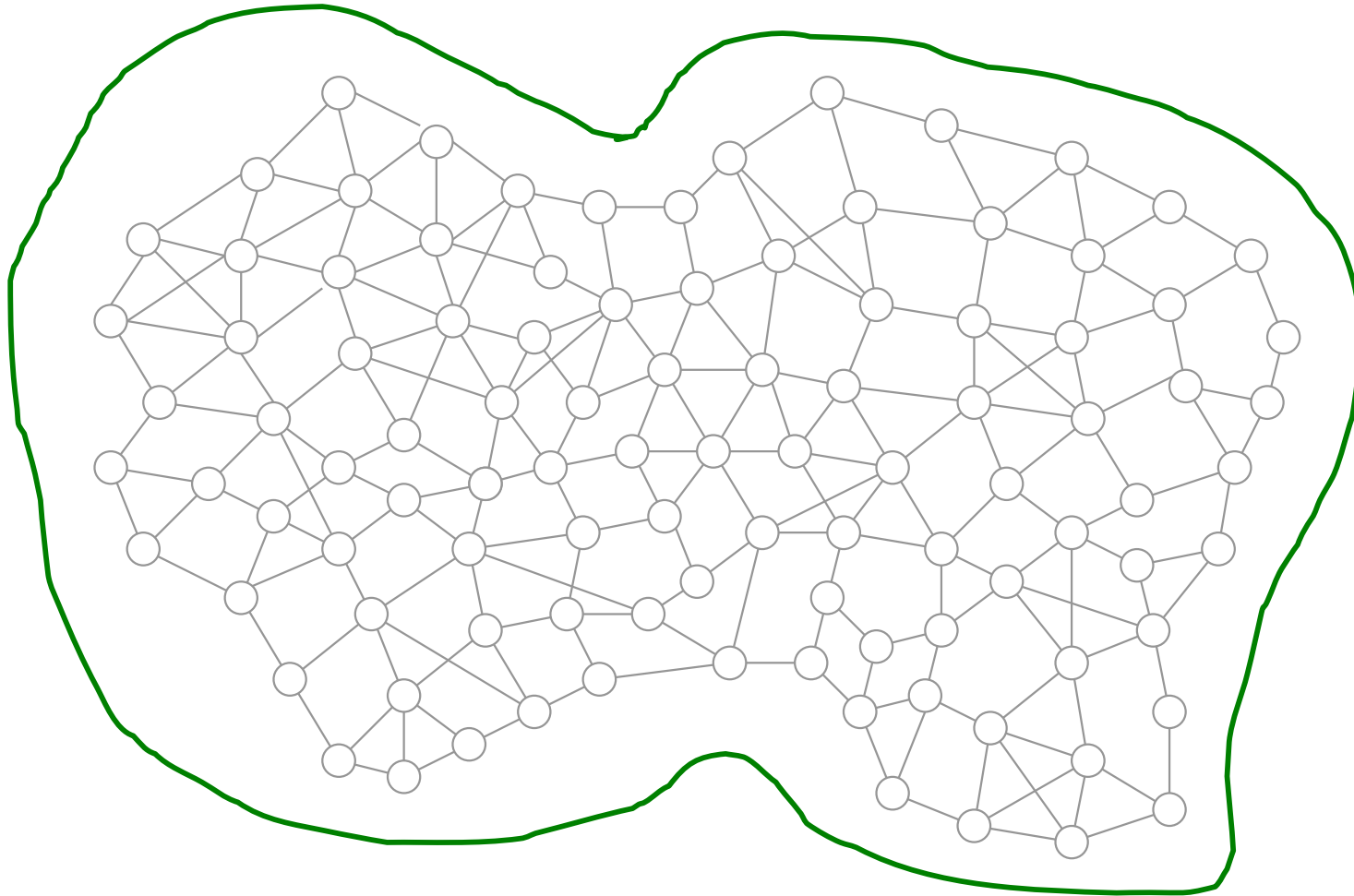
Protocol cost: Protocol's path distance from requesting node to object node

Existing Approaches

Protocol	Update Stretch	Network Kind	Runs on
Arrow [DISC'98]	$O(S_{ST})=O(D)$	General	Spanning tree
Relay [OPODIS'09]	$O(S_{ST})=O(D)$	General	Spanning tree
Combine [SSS'10]	$O(S_{OT})=O(D)$	General	Overlay tree
Ballistic [DISC'05]	$O(\log D)$	Constant-doubling	Hierarchical clustering with independent sets
Spiral [IPDPS'12]	$O(\log^2 n \log D)$	General	Hierarchical clusters

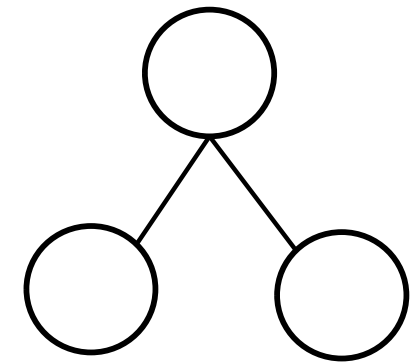
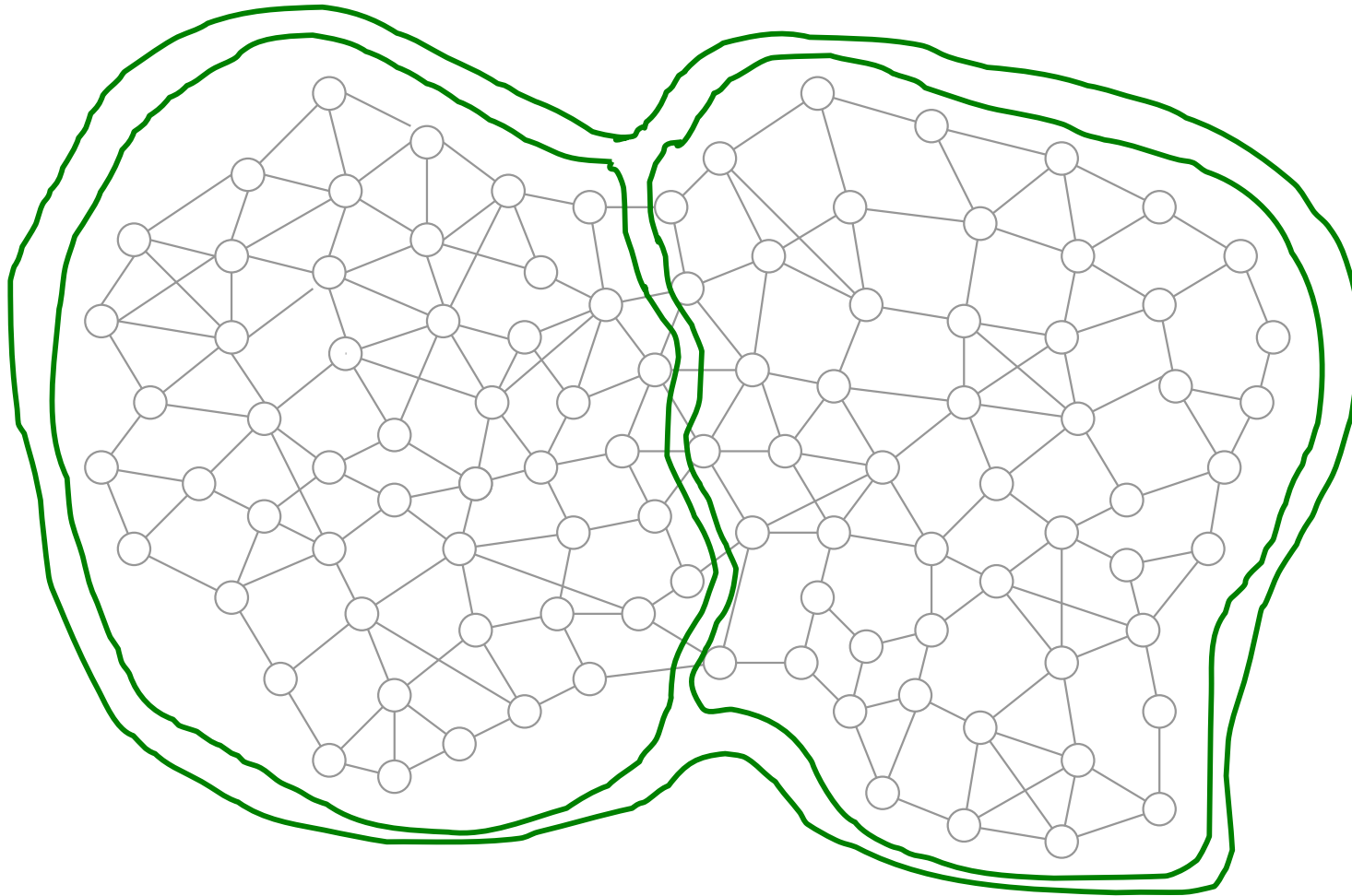
Stretch: Ratio of Protocol cost over Optimal cost

Approach: Hierarchical



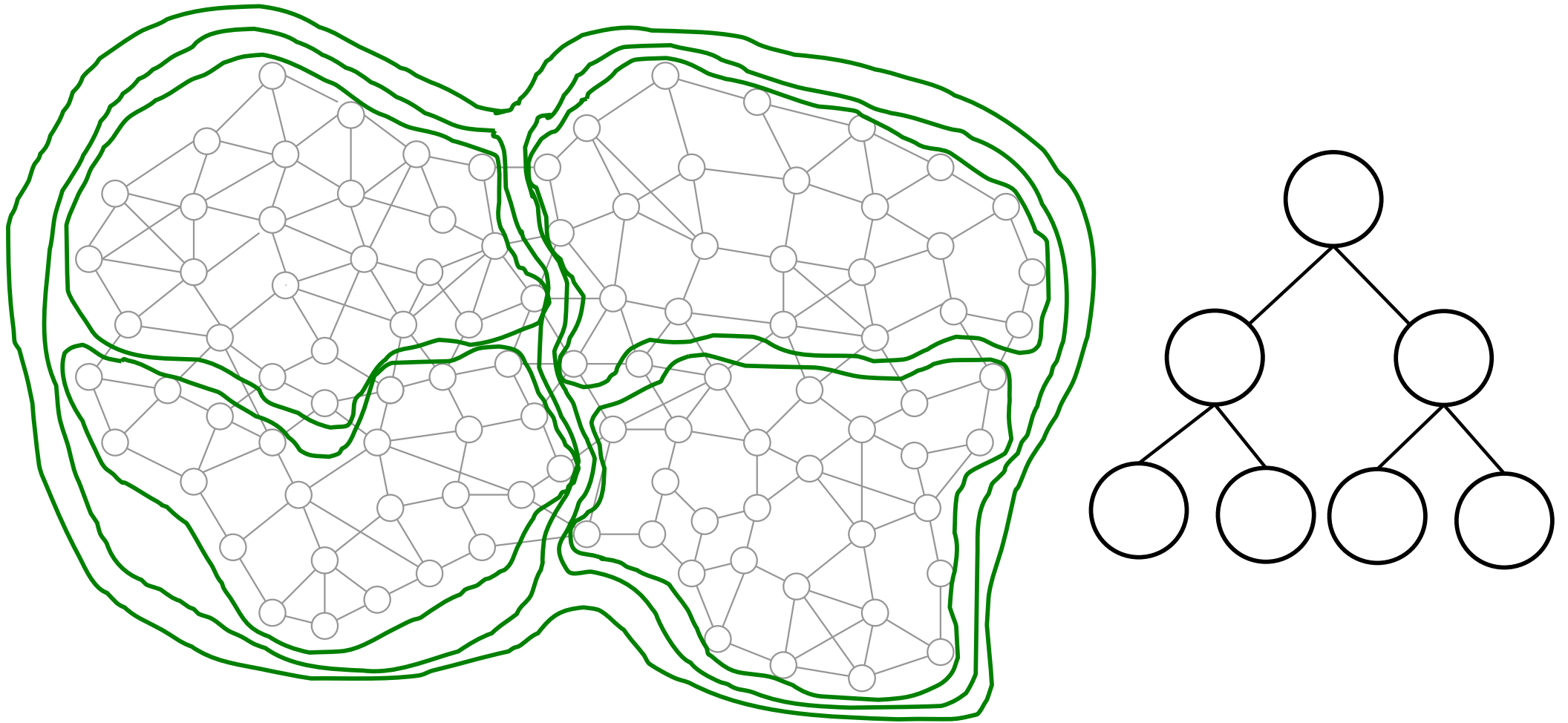
Network graph

Approach: Hierarchical



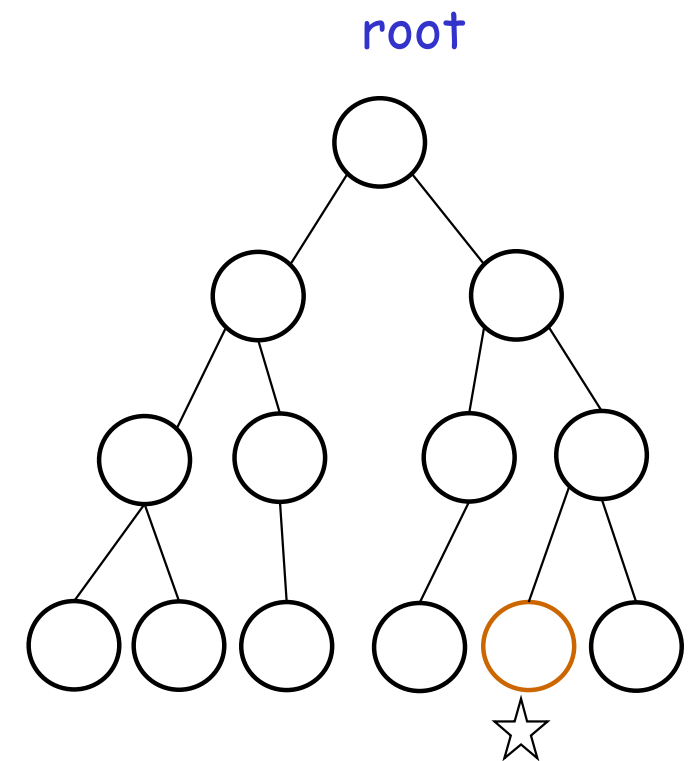
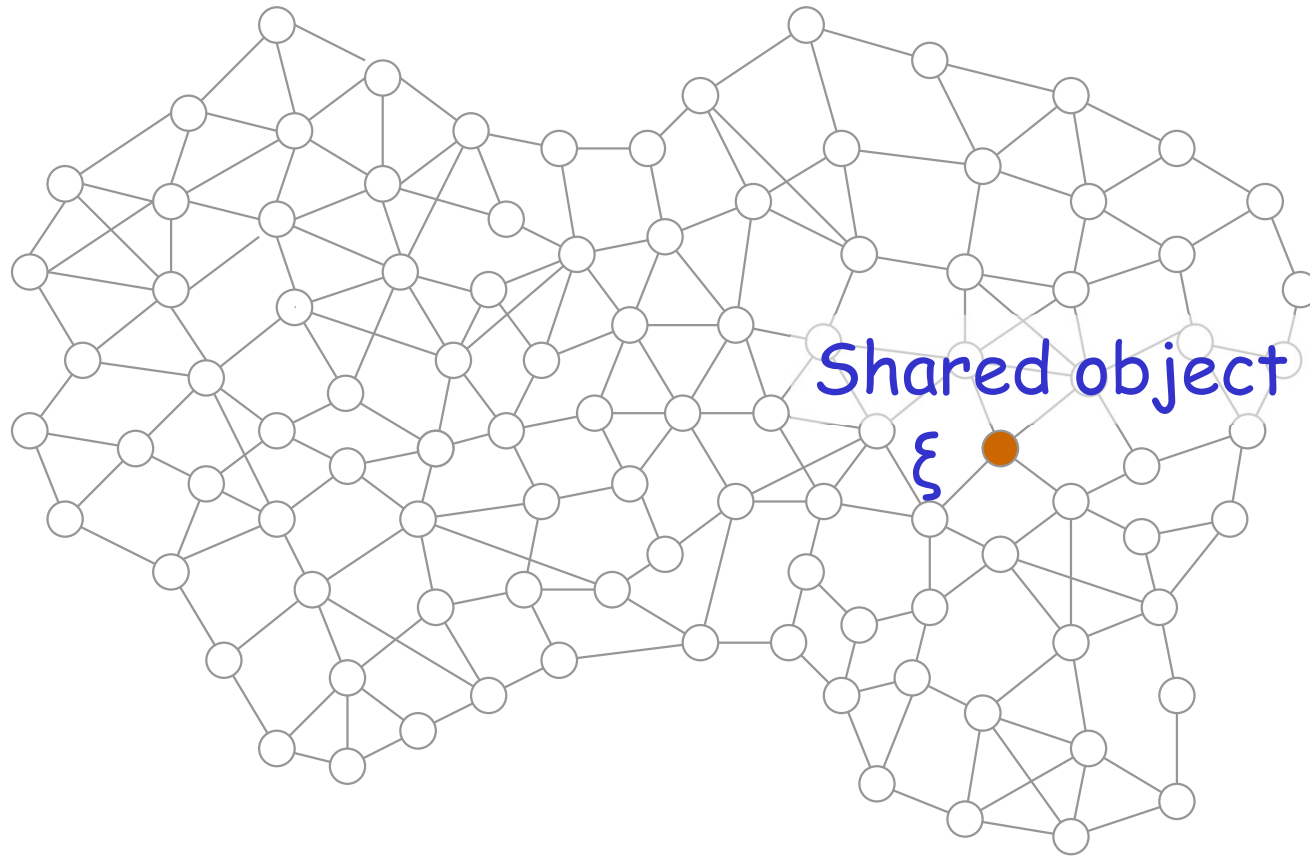
Alternative
representation as a
hierarchy tree with
leader nodes

At the lowest level (level 0) every node is a cluster

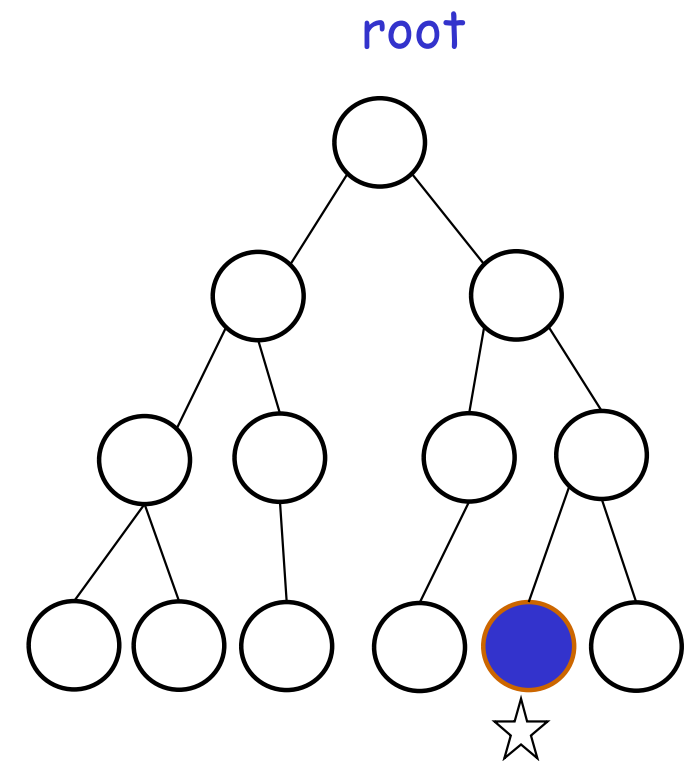
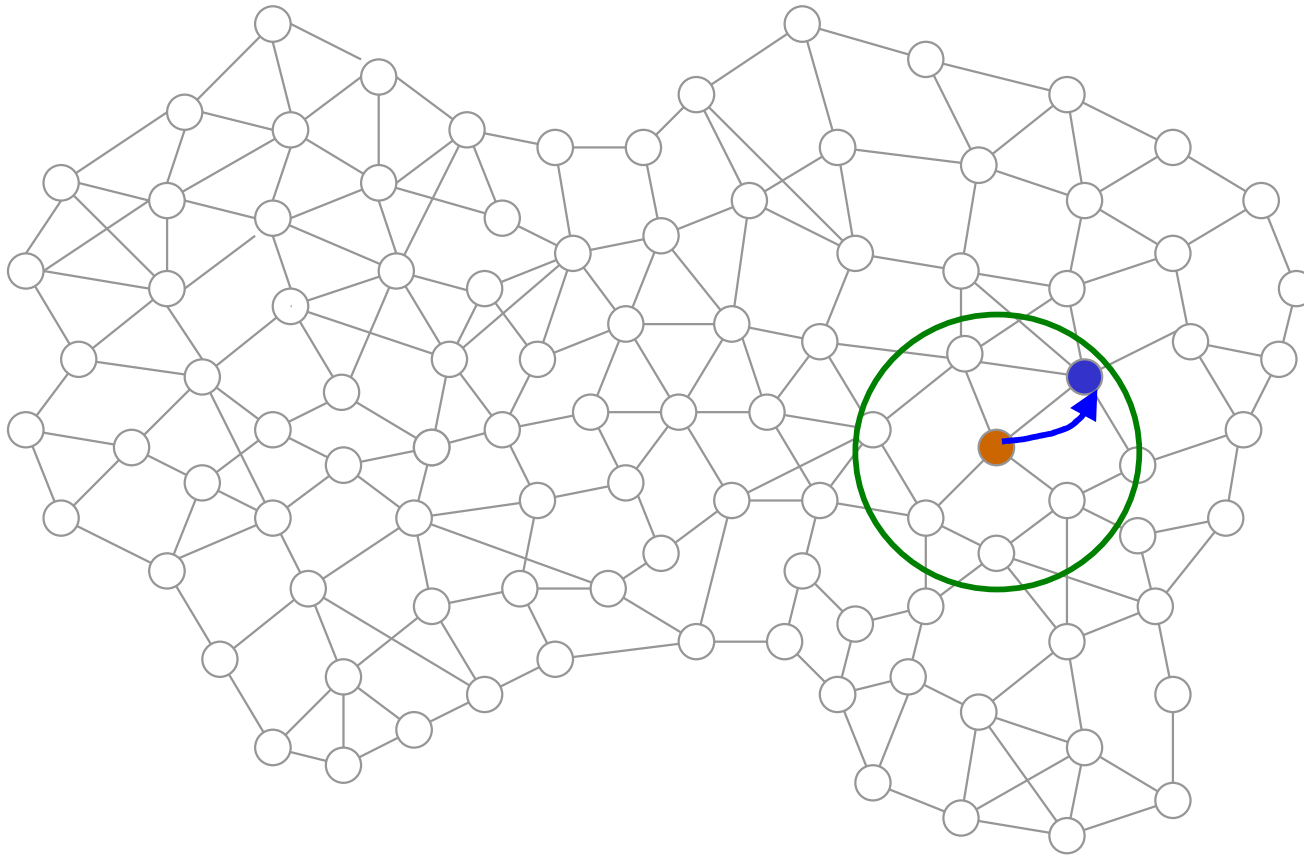


Directories at each level cluster, downward pointer if object locality known

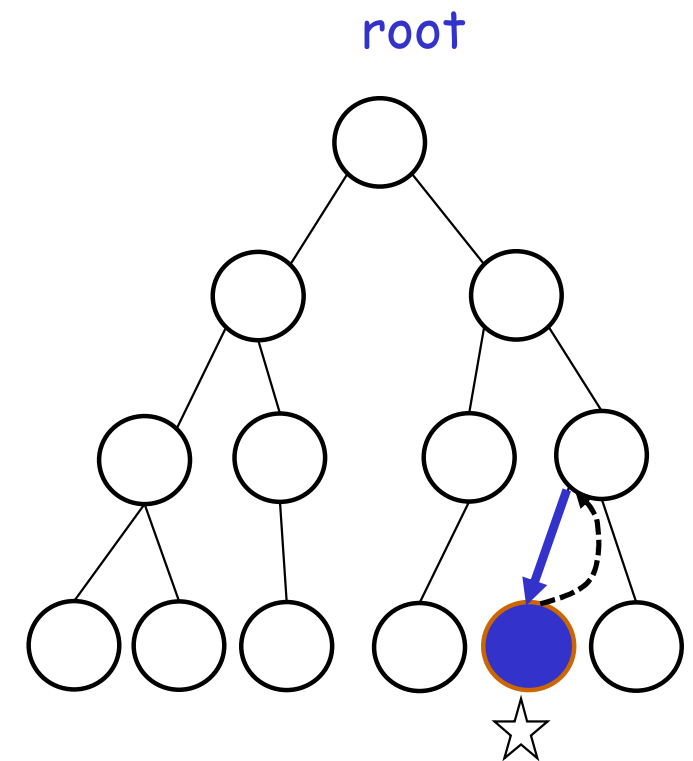
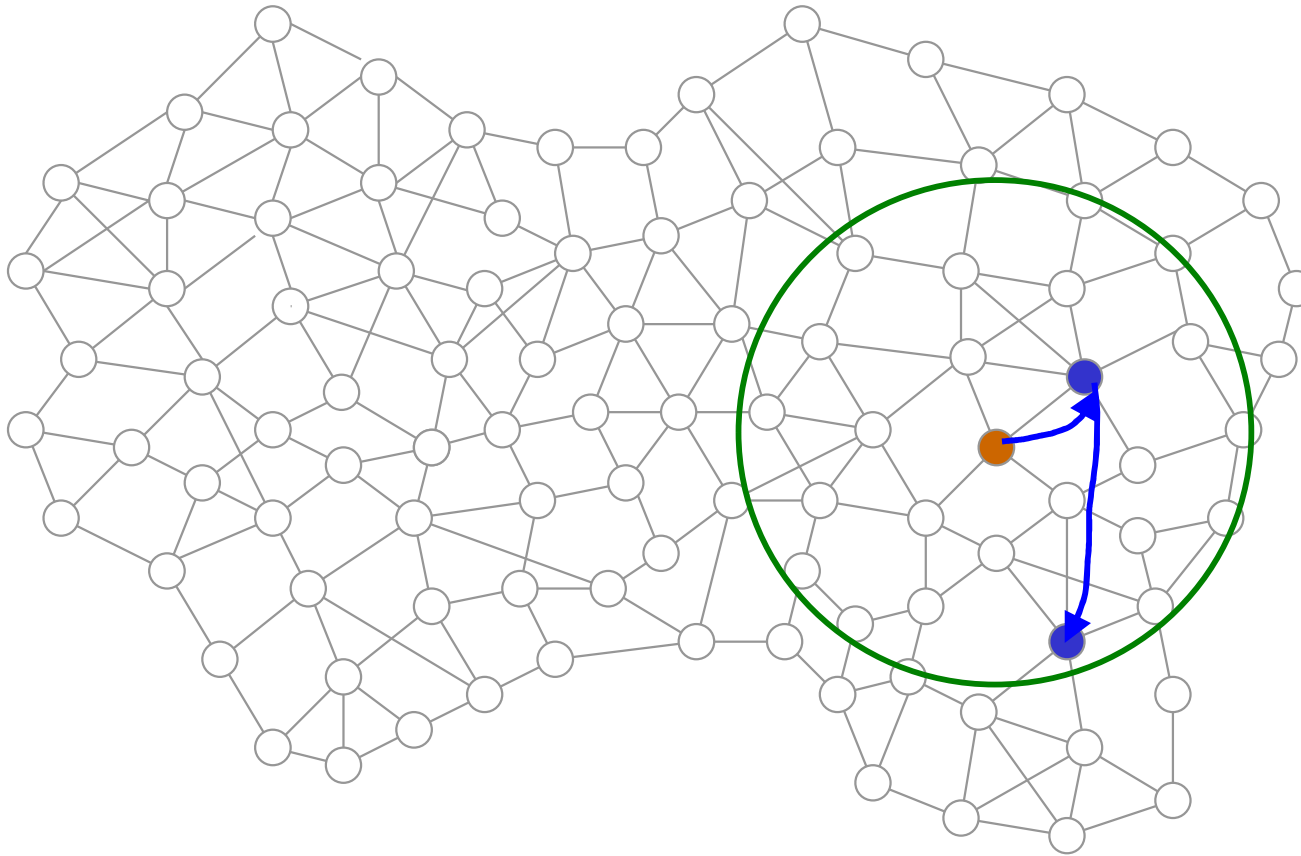
Publish object



Send request to the leader

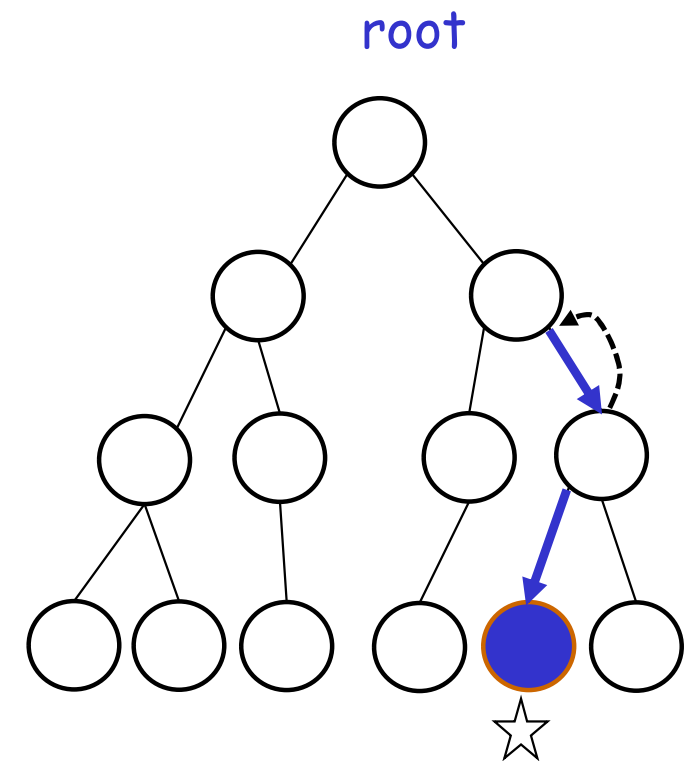
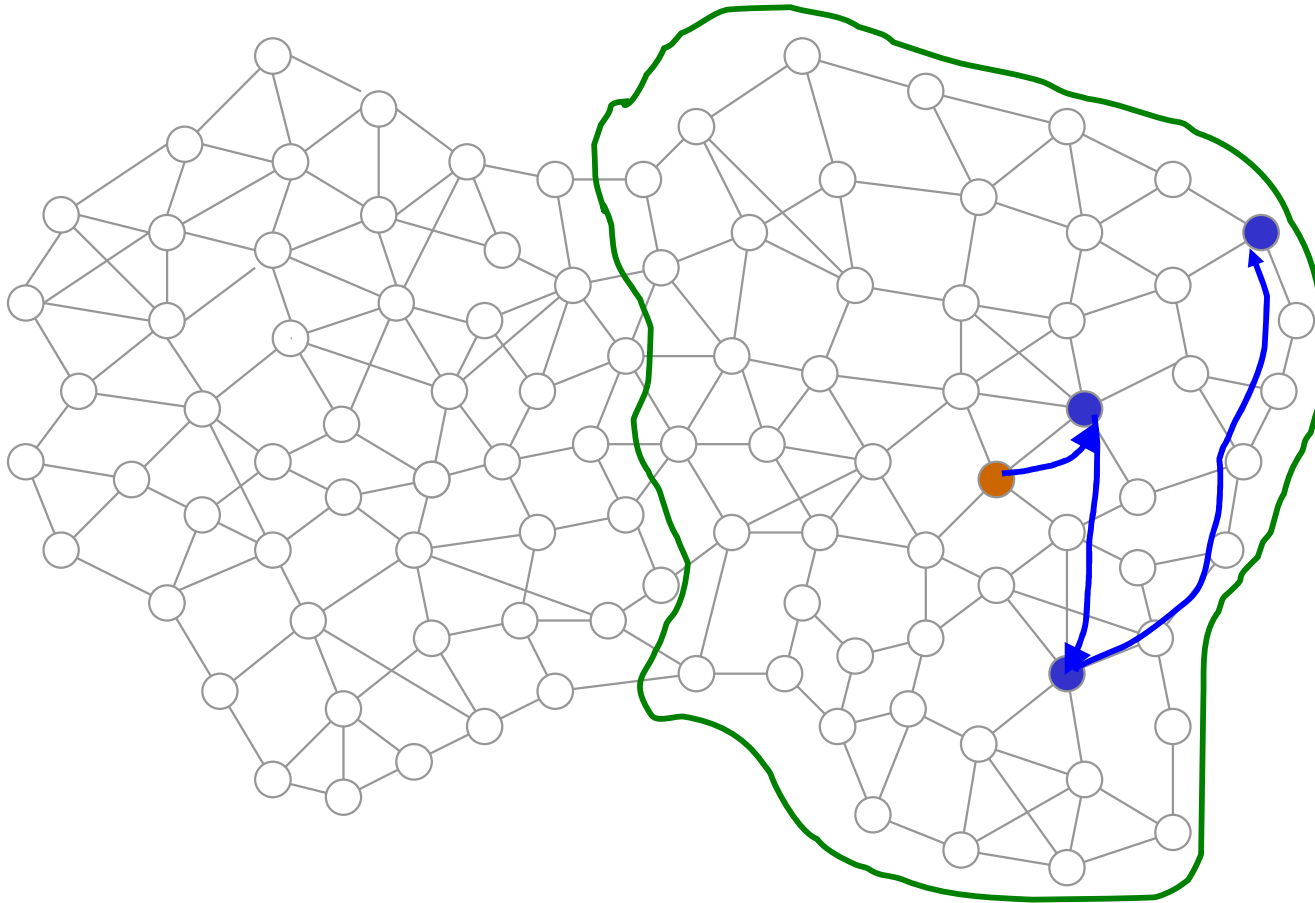


Continue up phase



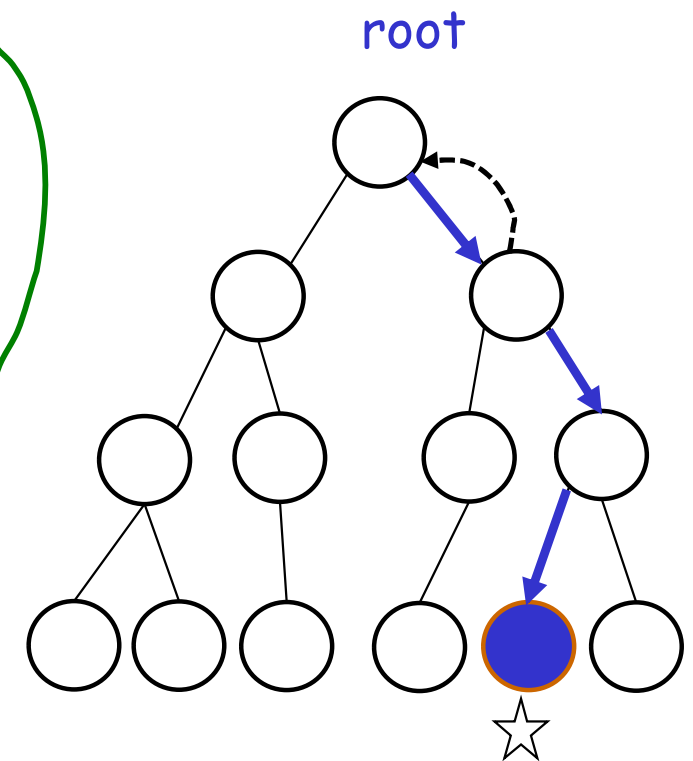
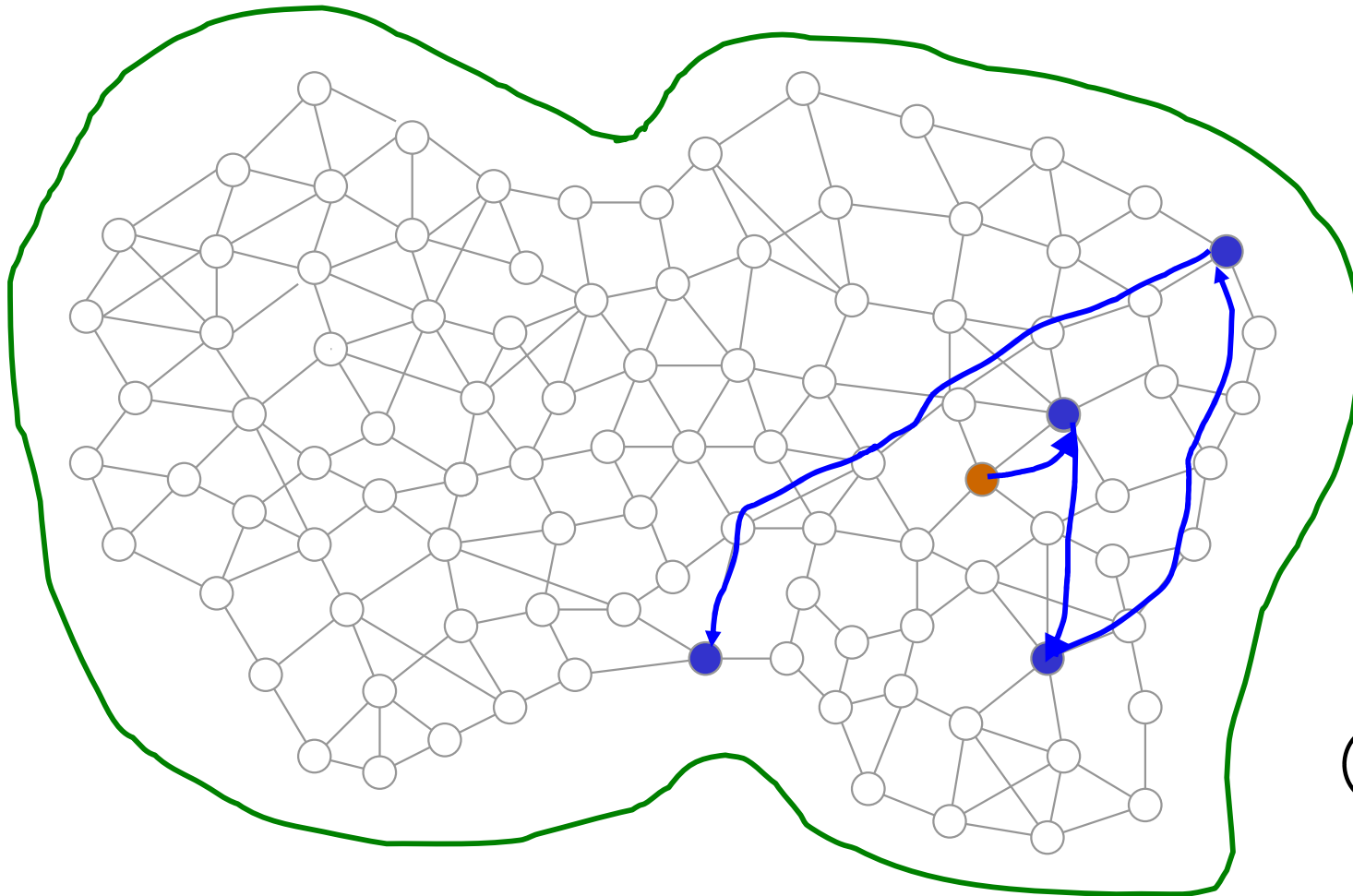
Set downward pointer while going up

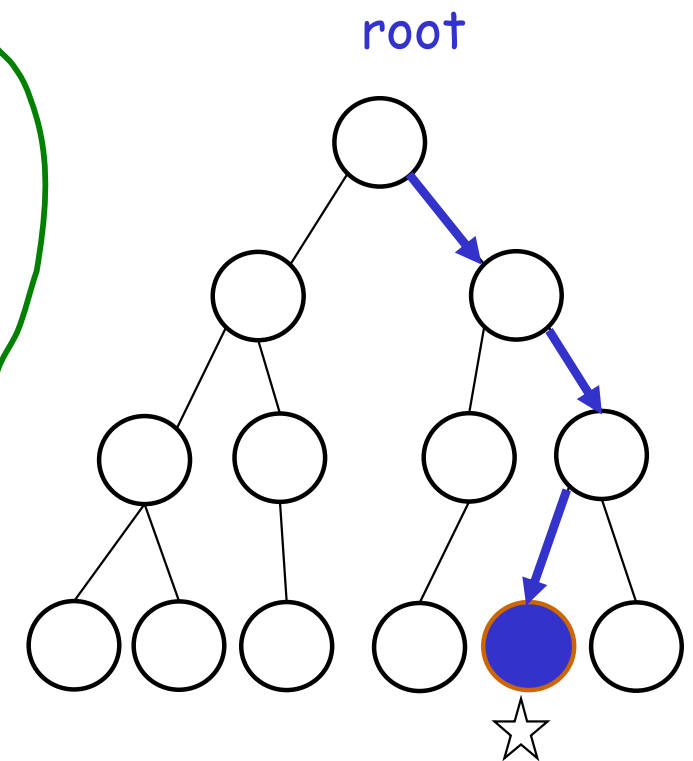
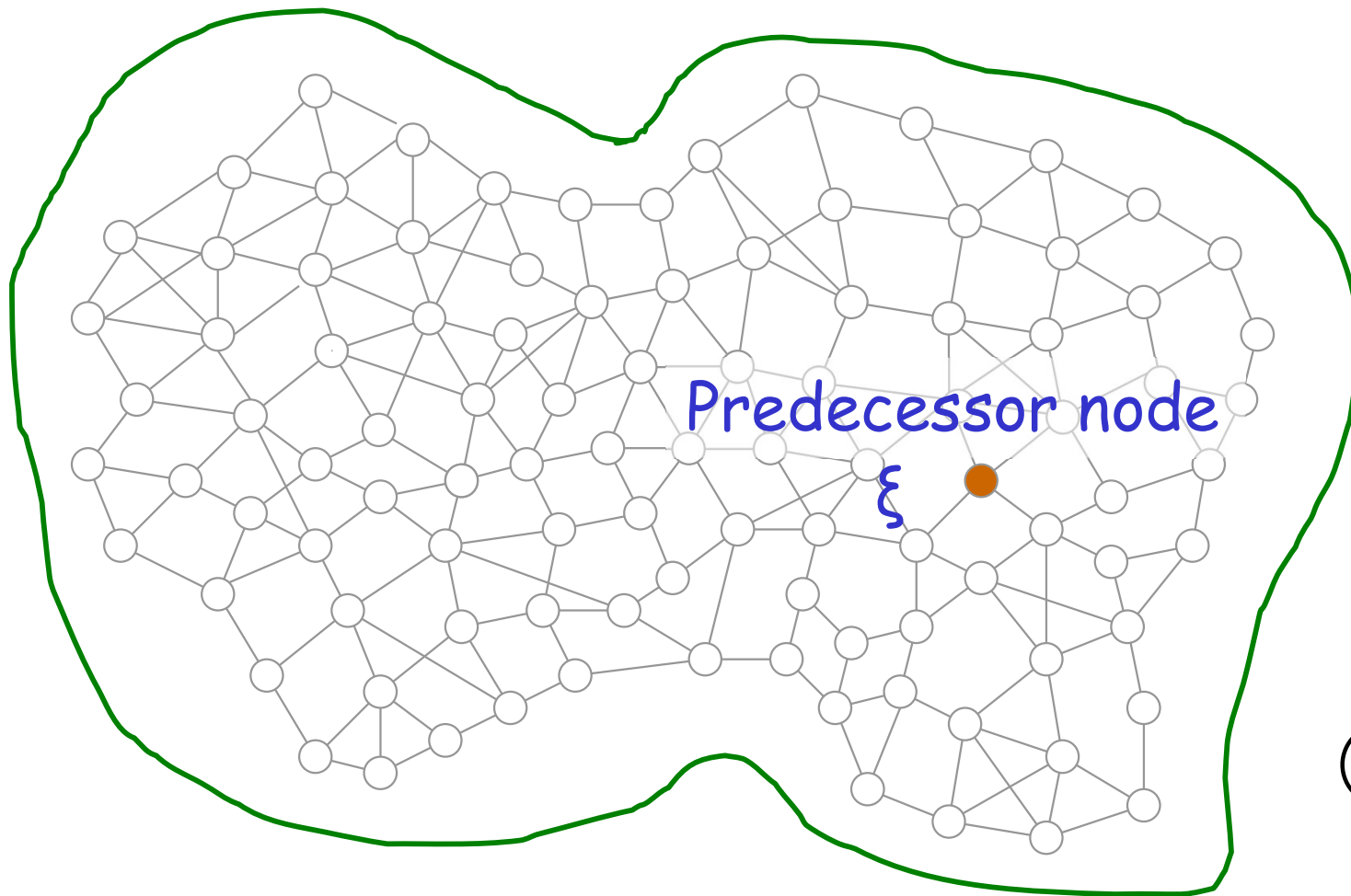
Continue up phase



Set downward pointer while going up

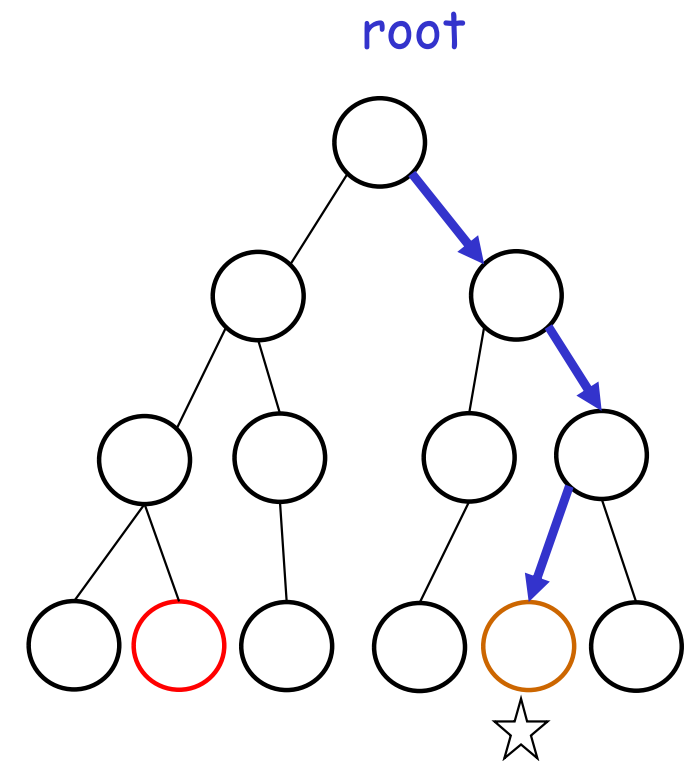
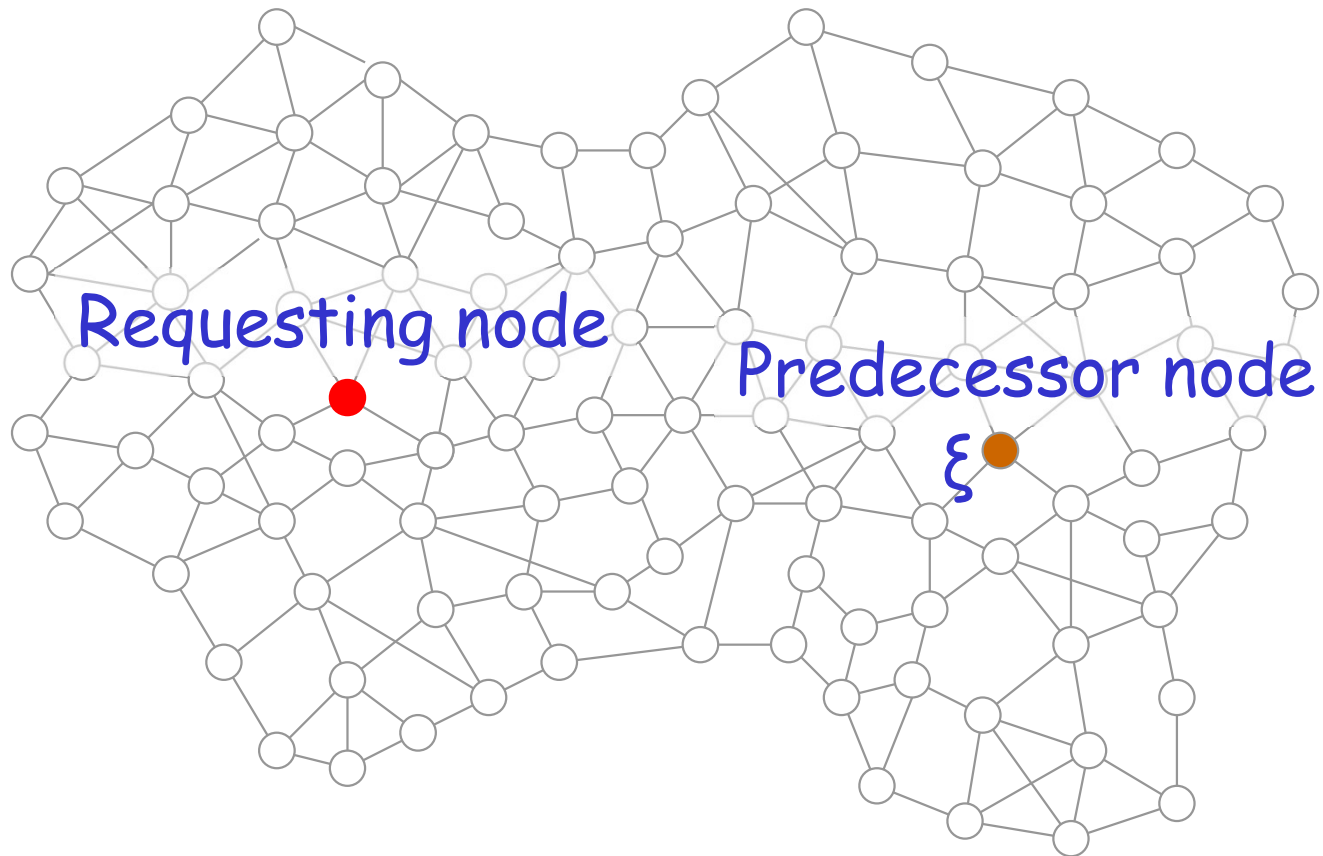
Root node found, stop up phase



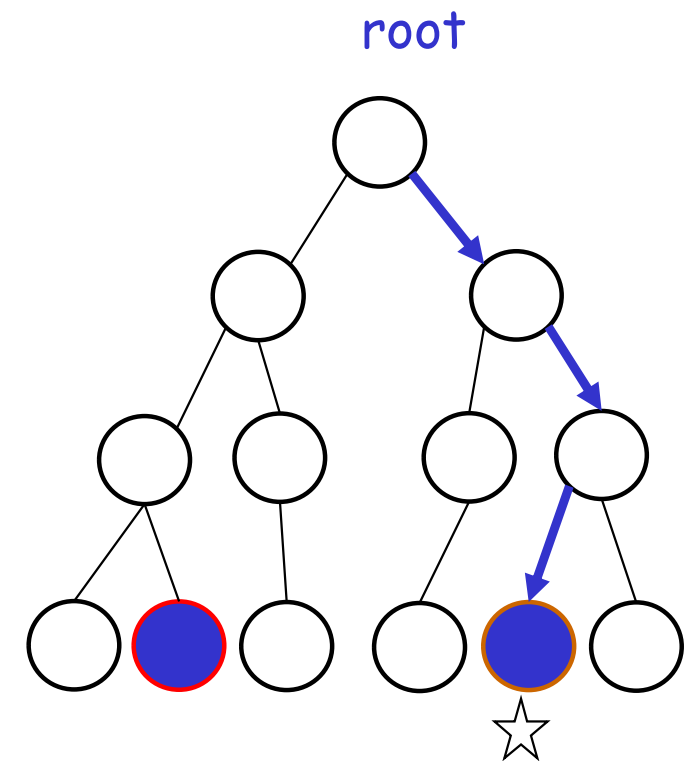
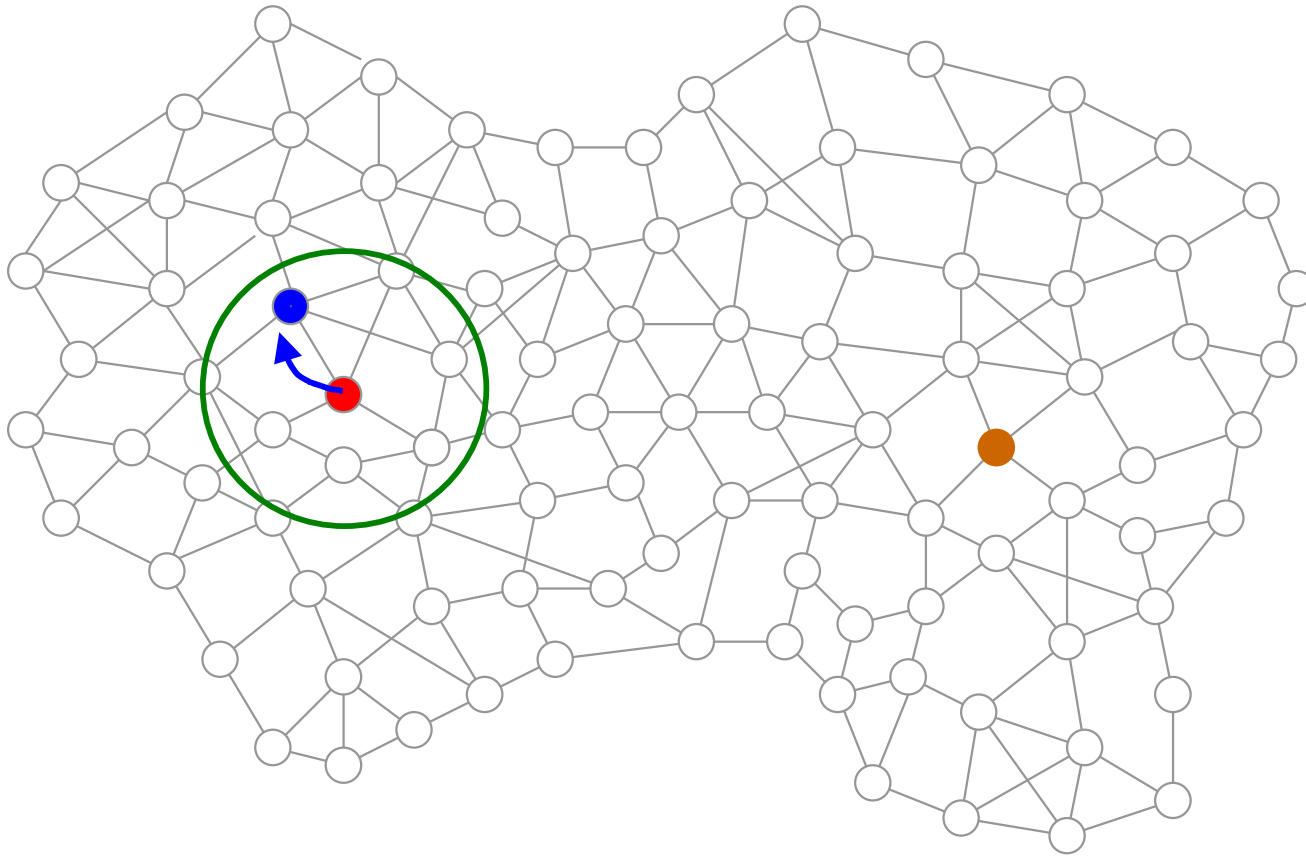


A successful Publish operation

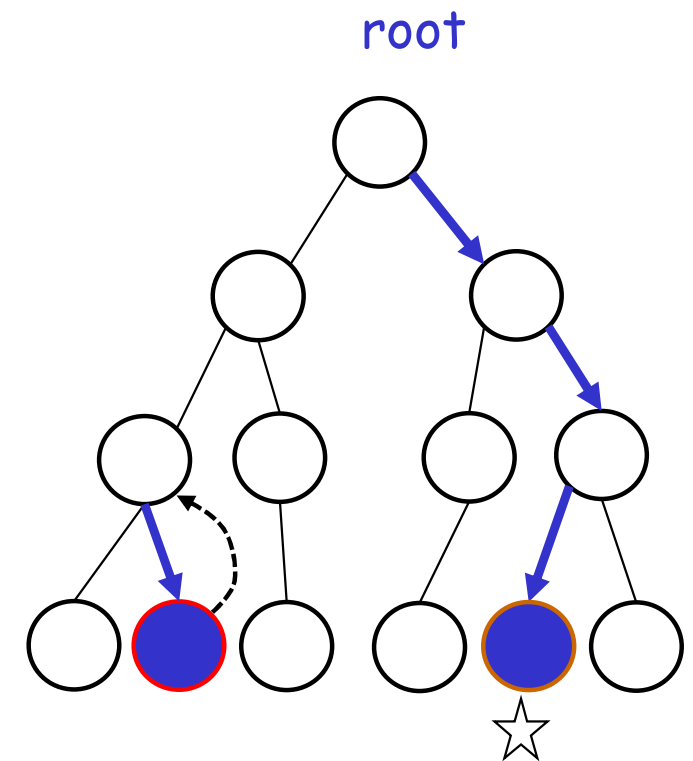
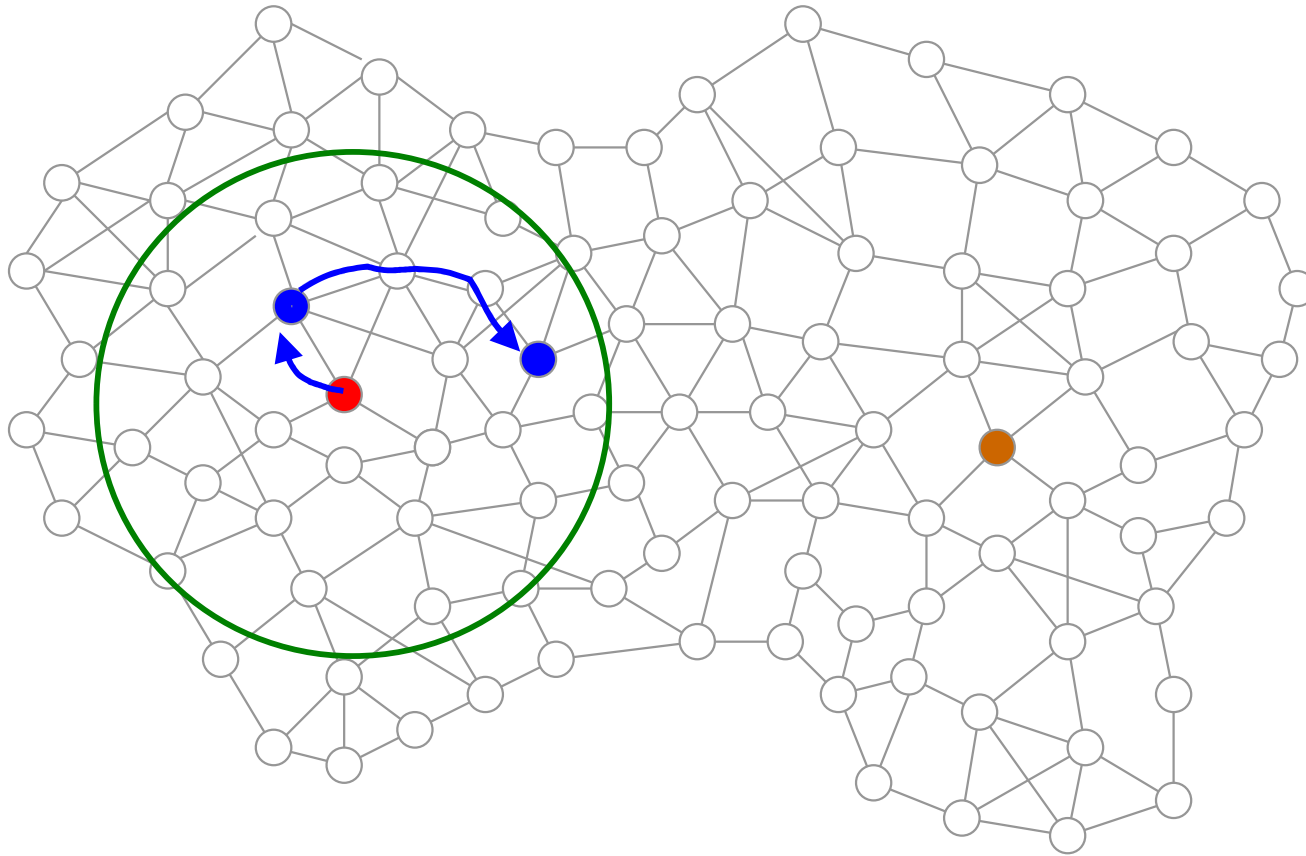
Move object (for Write)



Send request to leader node of the cluster upward

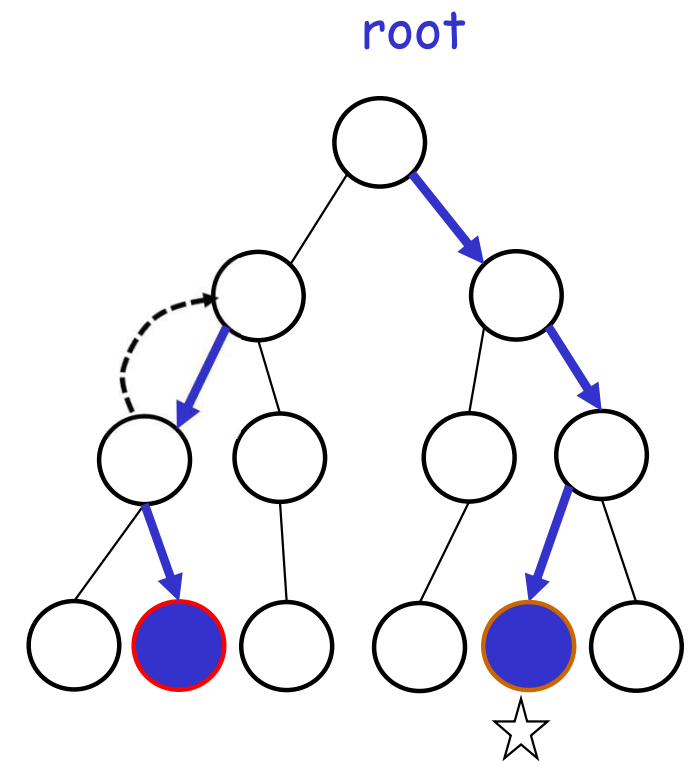
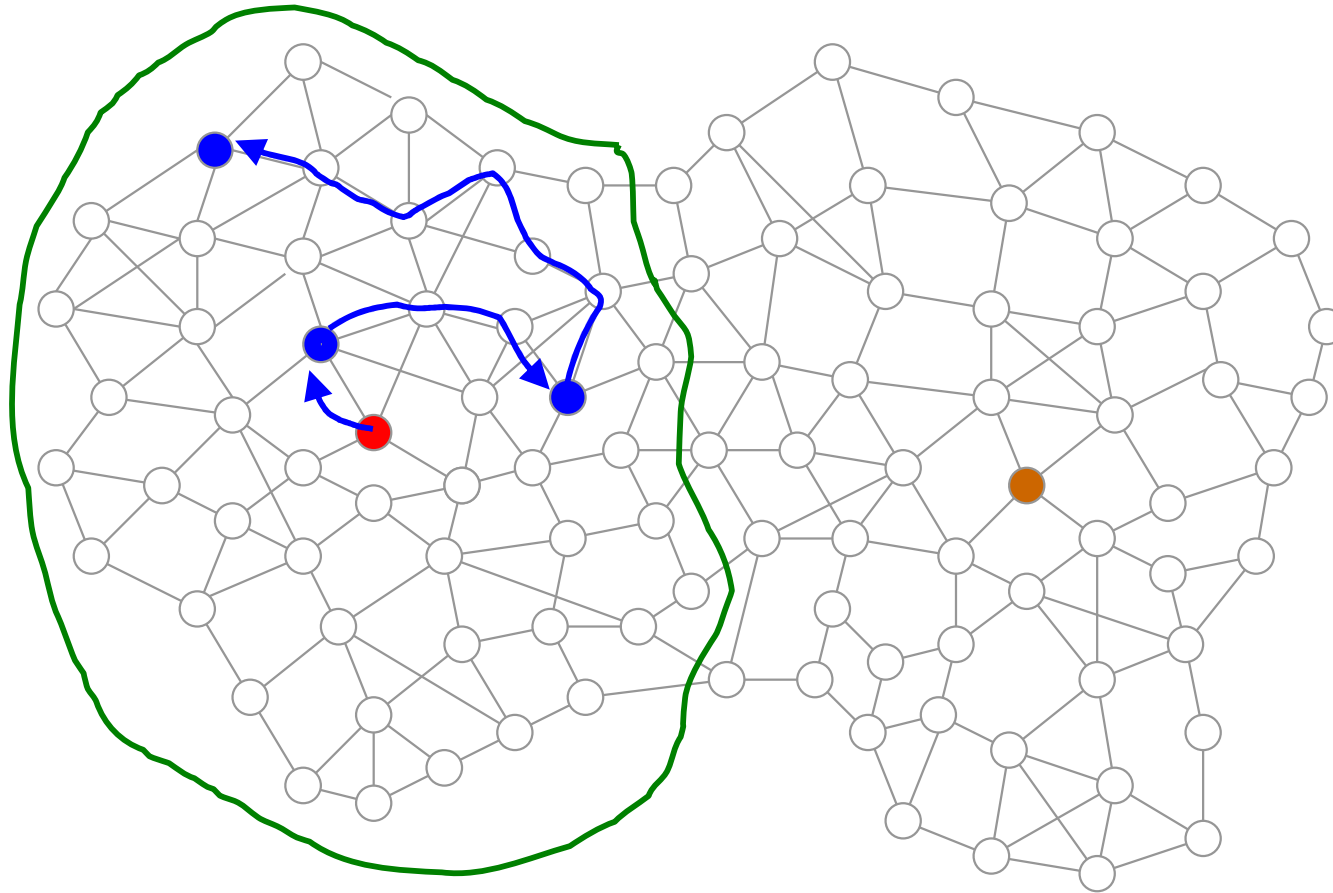


Continue up phase until downward pointer found



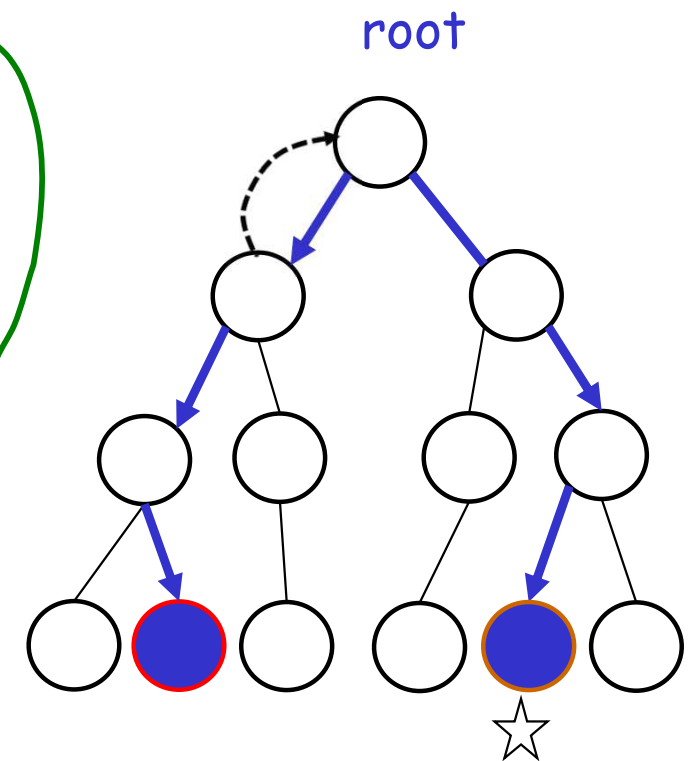
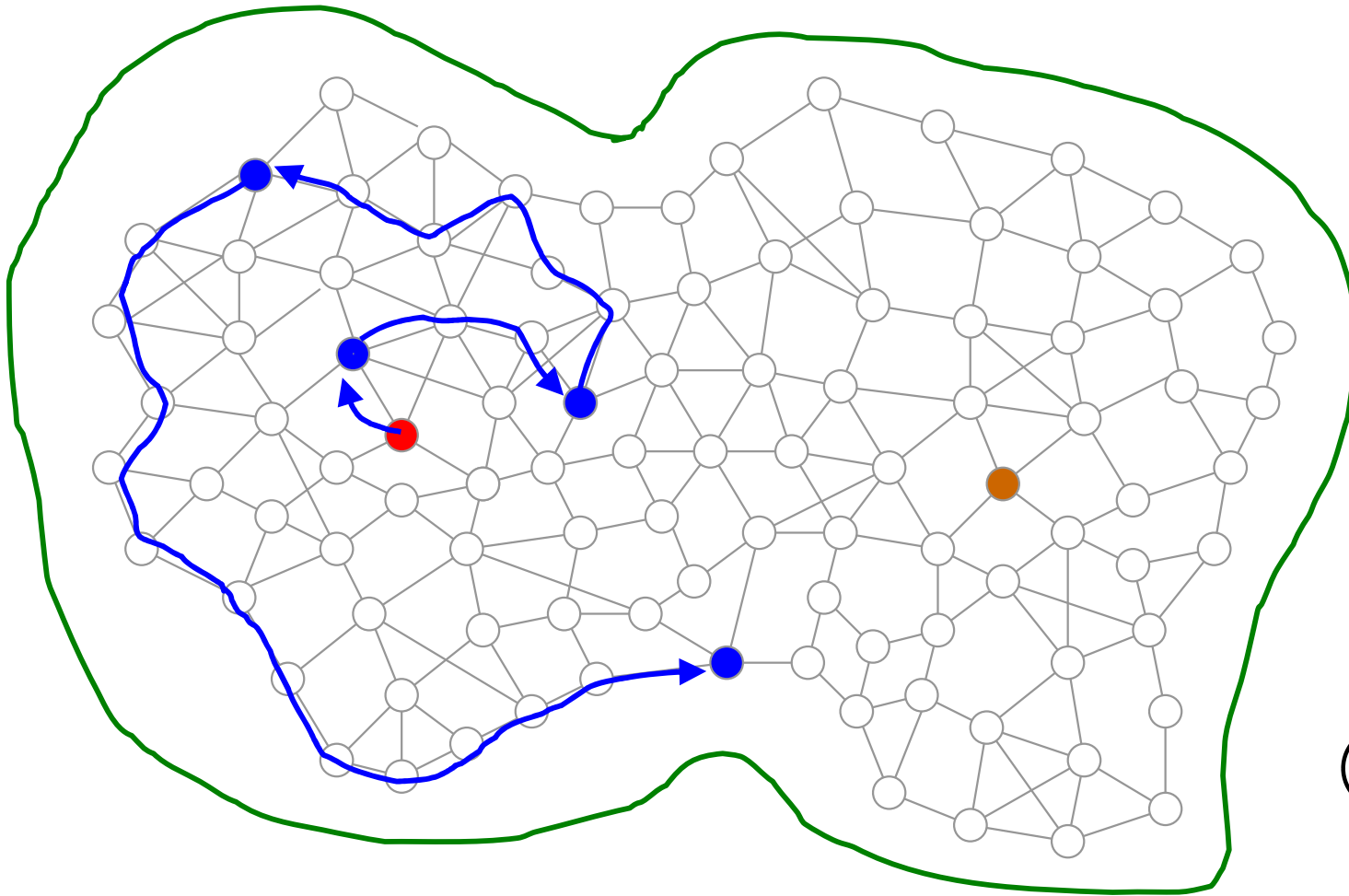
Set downward pointer while going up

Continue up phase



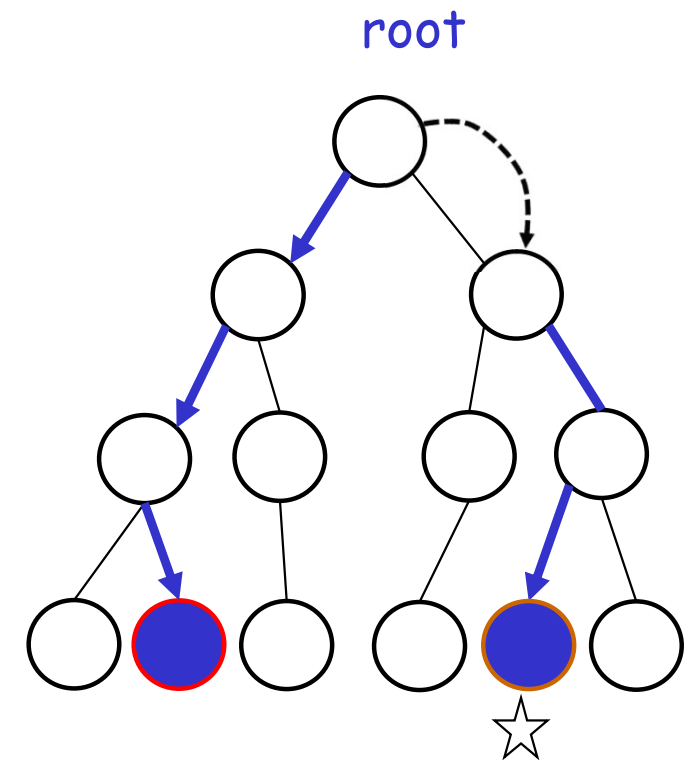
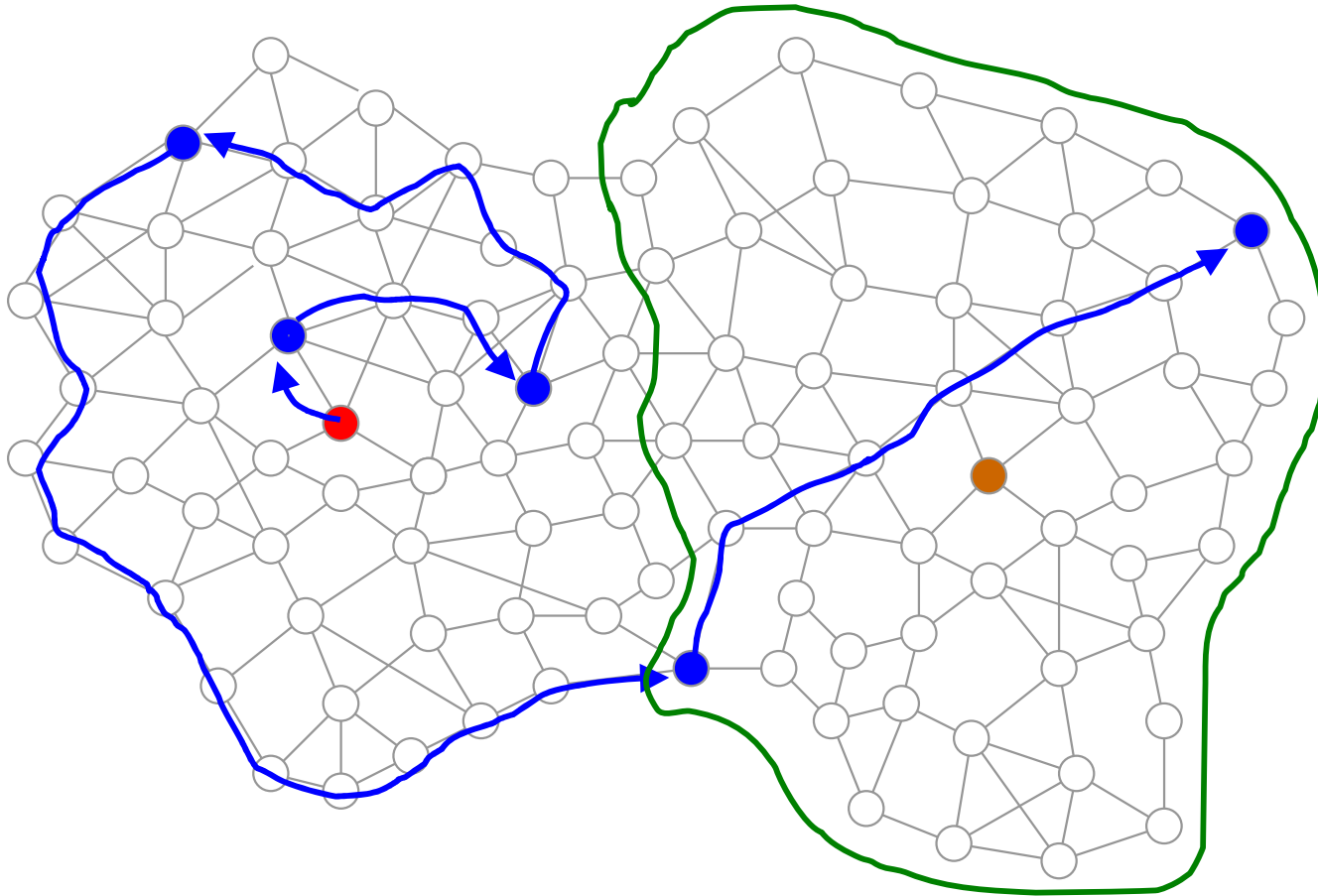
Set downward pointer while going up

Continue up phase



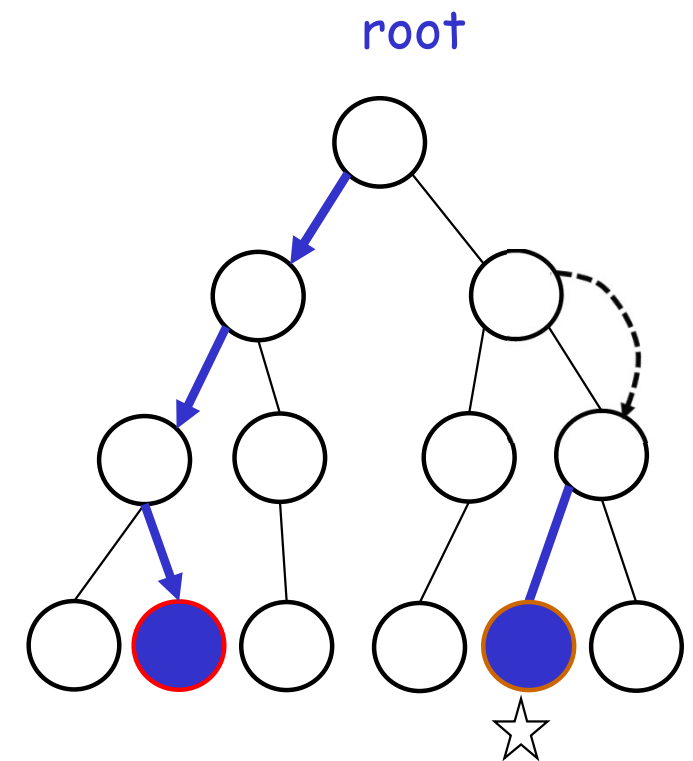
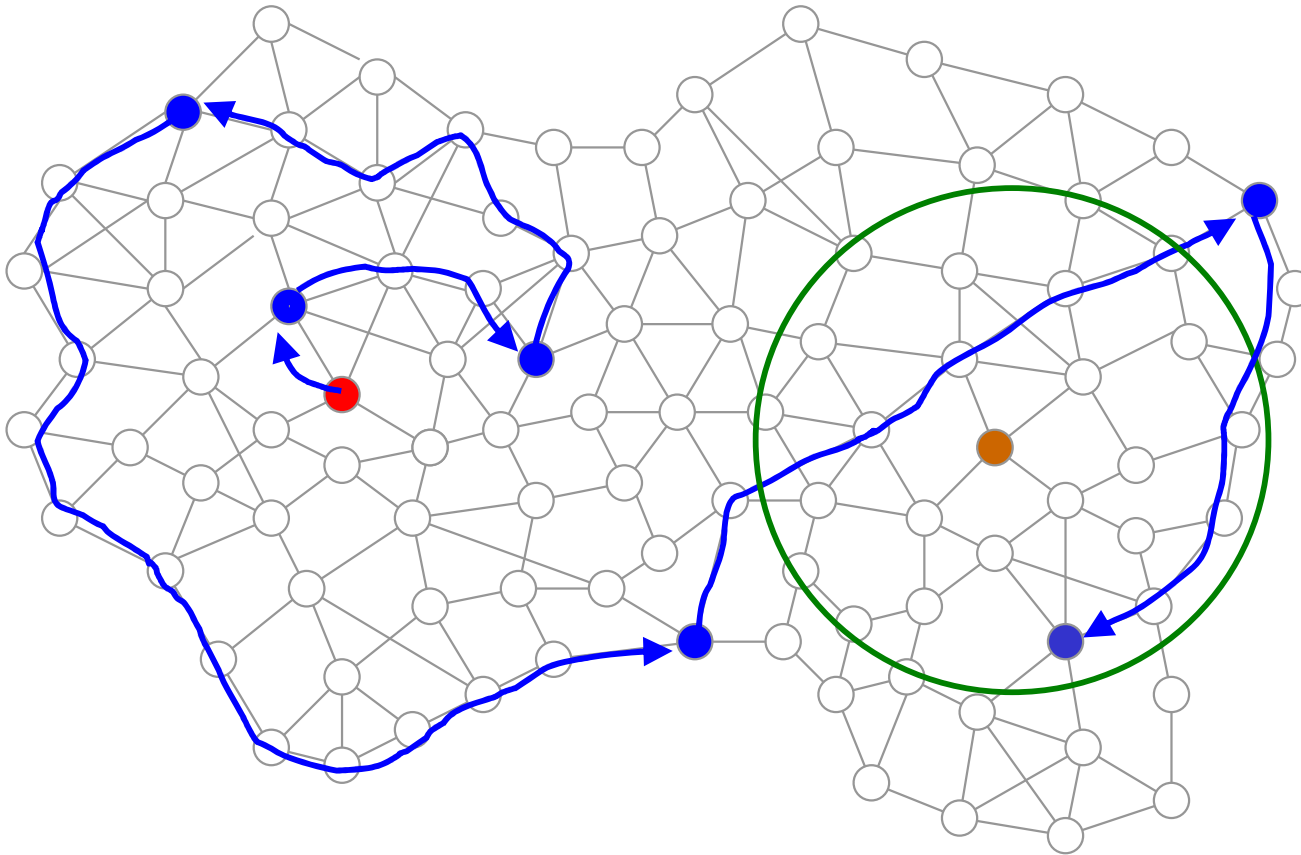
Set downward pointer while going up

Downward pointer found, start down phase



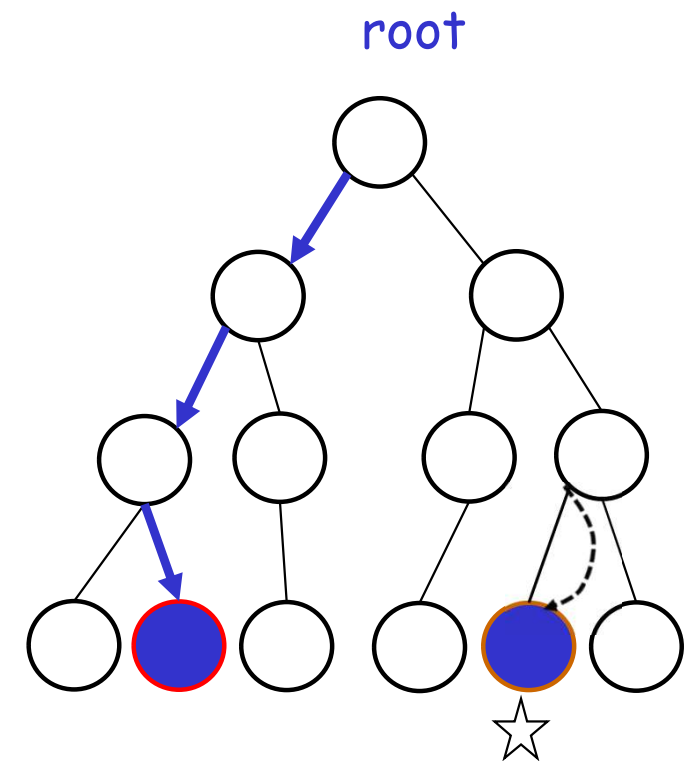
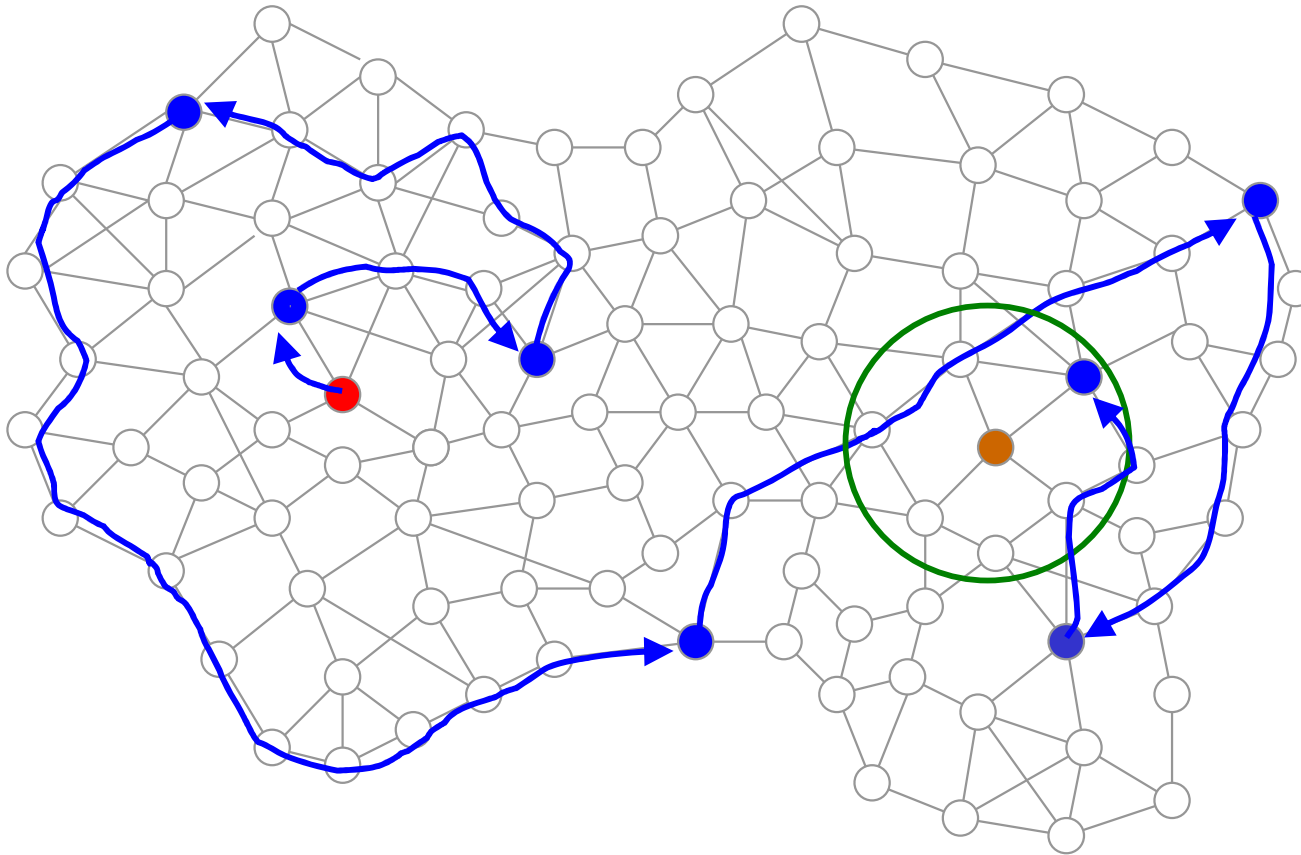
Discards downward pointer while going down

Continue down phase



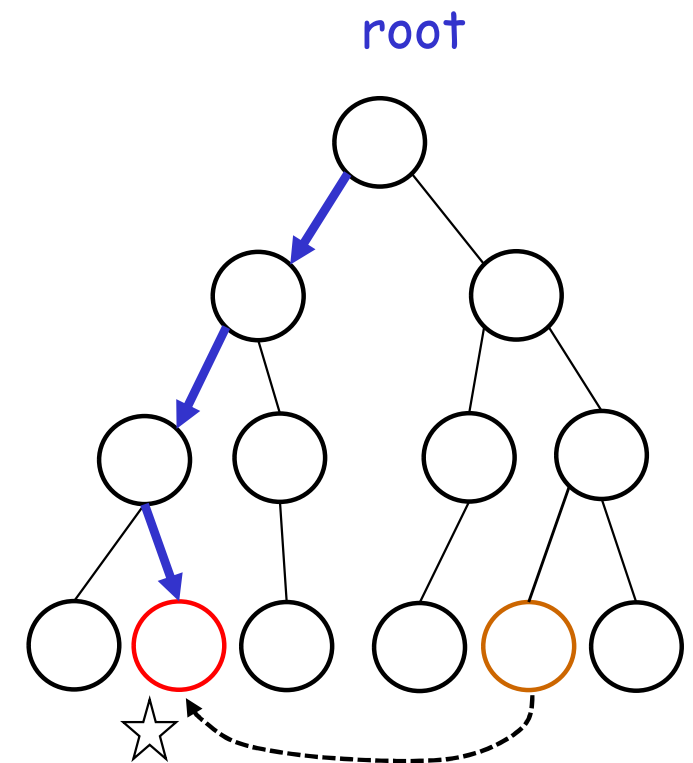
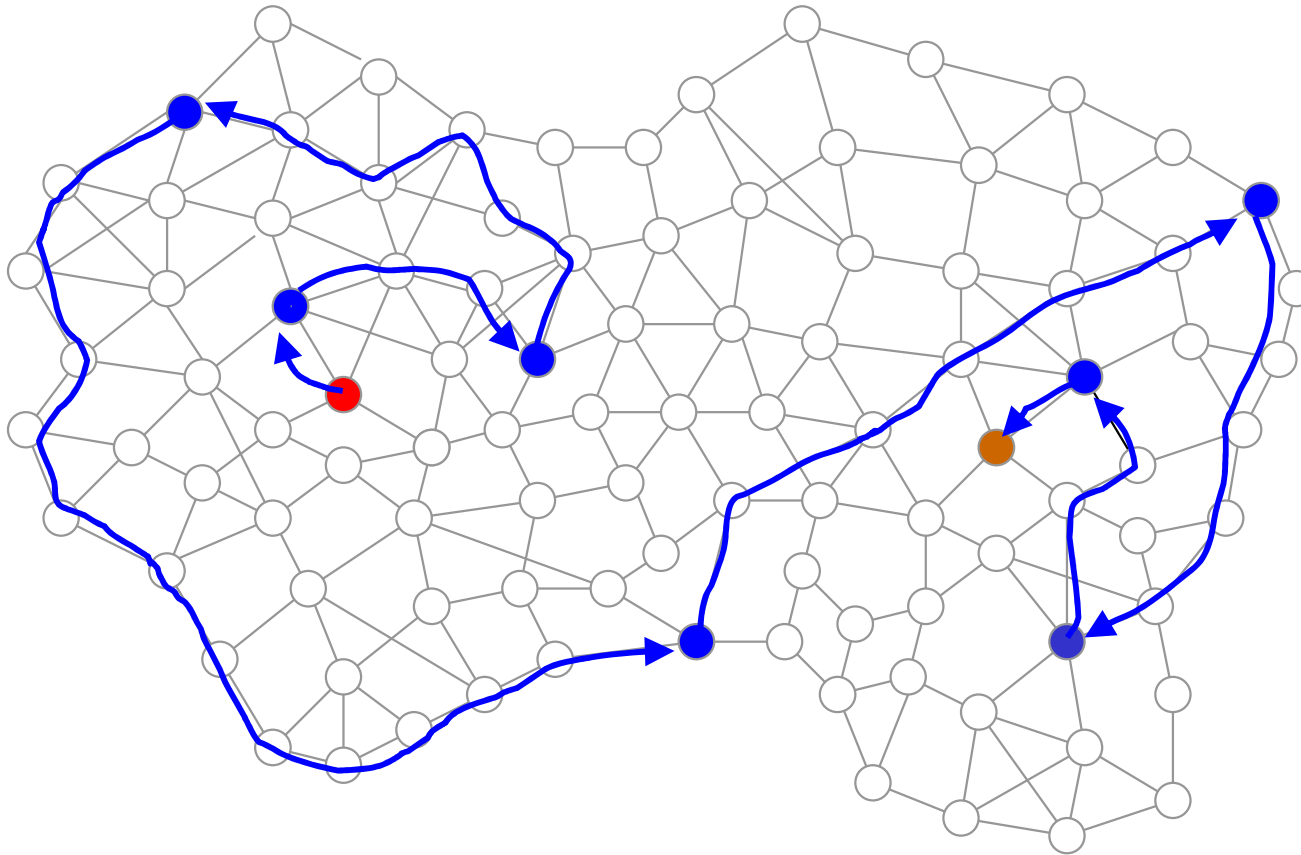
Discards downward pointer while going down

Continue down phase



Discards downward pointer while going down

Predecessor reached, object is found and can be moved



Lookup object (for Read) is similar

Landscape of Existing Research

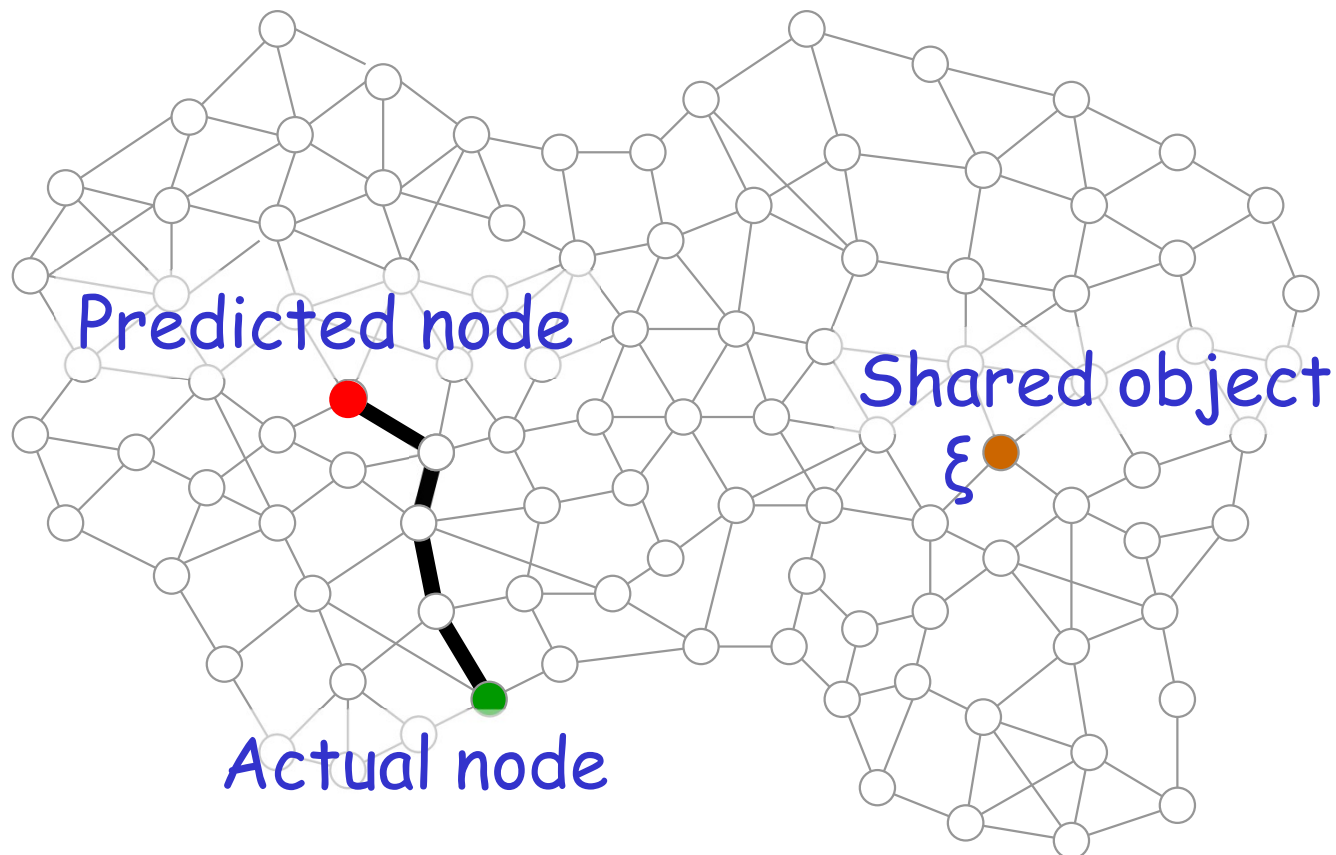
- **Online Model - Move/Lookup (write/read) requests come over time**
 - Requesting nodes not known a priori
 - Request arrival times not known a priori

New Line of Research

- **Online Model - Move/Lookup (write/read) requests come over time**
 - Requesting nodes not known a priori
 - Request arrival times not known a priori
- **Prediction Model - Can the protocols be optimized (w.r.t. stretch) if**
 - Some limited knowledge on input is provided as predictions

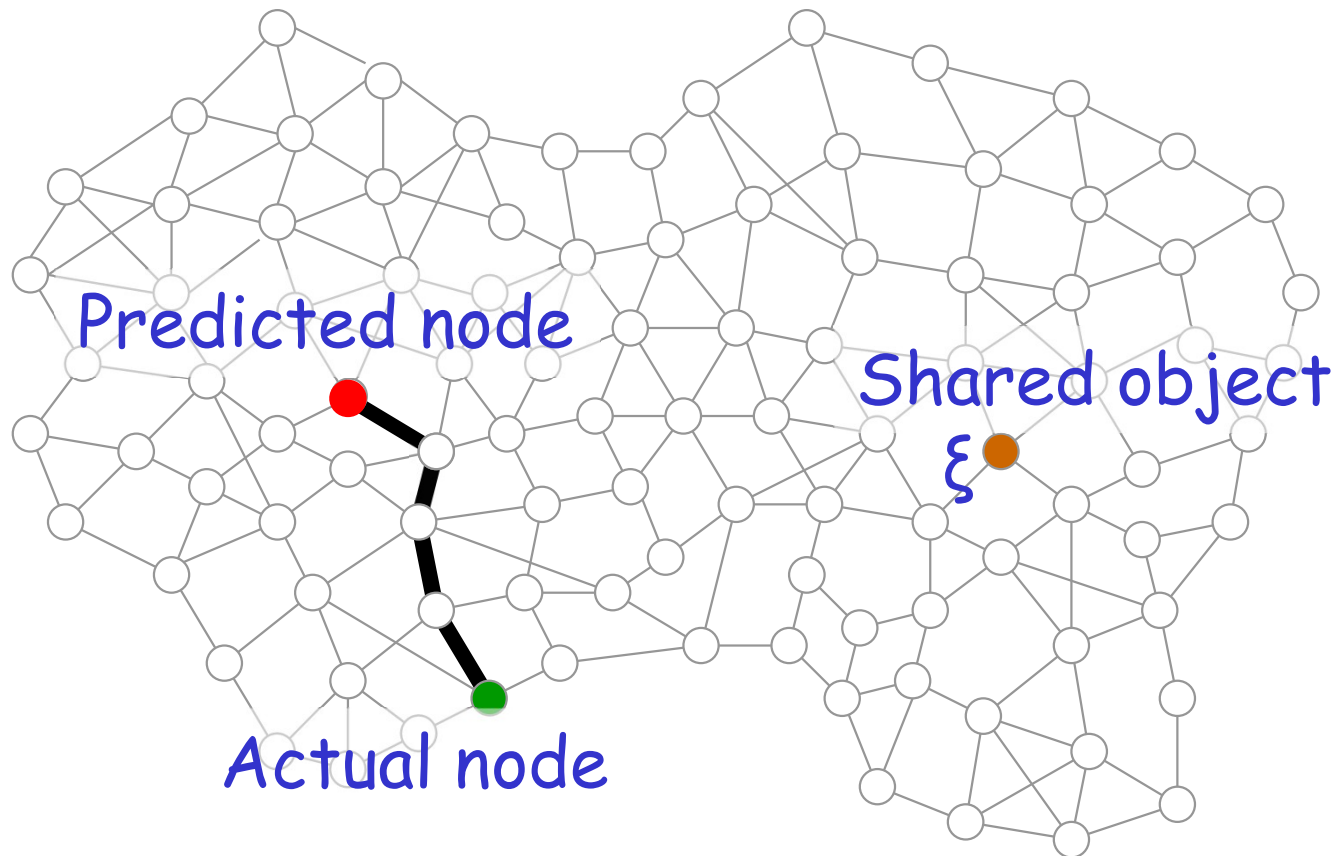
Location Prediction Model

- Requesting nodes are *predicted* a priori
- Request arrival times not known a priori
- Predictions can be erroneous (not known both error and its amount)



Location Prediction Model

- **Prediction error:** shortest distance between predicted node and actual node
- **Maximum Prediction Error:** D , the diameter



Protocol Design in the Prediction Model

- **Consistency**
 - Stretch when no error (predictions are perfect)
- **Smoothness**
 - Stretch when predictions have error
- **Robustness**
 - Stretch when predictions have maximum error

Protocol Design in the Prediction Model

- **Consistency**
 - Stretch when no error (predictions are perfect)
 - **Goal: much better than online (no prediction) model**
- **Smoothness**
 - Stretch when predictions have error
 - **Goal: gradually degrading with increasing error**
- **Robustness**
 - Stretch when predictions have maximum error (D in our case)
 - **Goal: no worse than online (no prediction) model**

New Results

Two Protocols in the Location Prediction Model

- **Tree-based**

- S_{steiner} -consistent, no error
- $(S_{\text{steiner}} + \text{Error})$ -smooth for error "Error"
- S_{ST} -robust, maximum error

- **Cluster-based**

- $O(\log^2 n \log D')$ -consistent, D' diameter of predicted nodes + initial owner node
- $O(\log^2 n \log(D' + \text{Error}))$, D' diameter of predicted nodes + initial owner node
- $O(\log^2 n \log D)$ -robust, maximum error

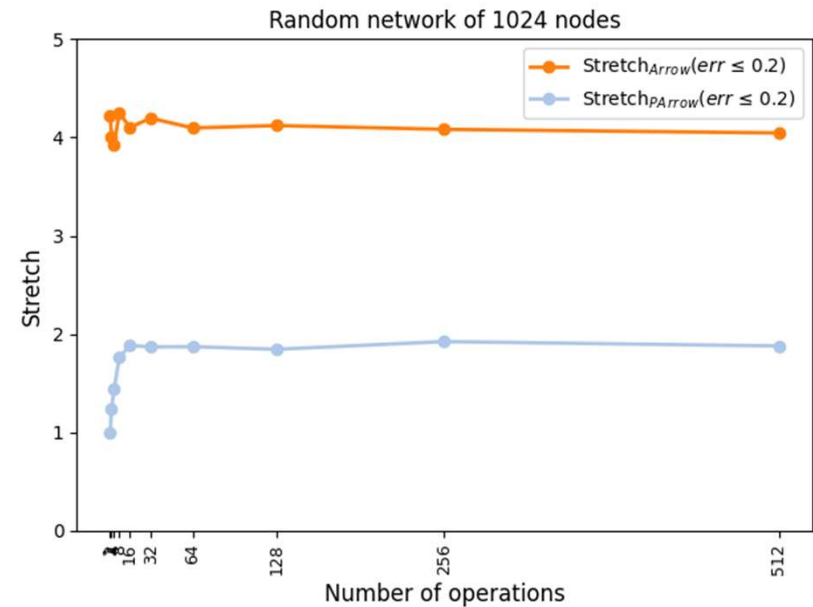
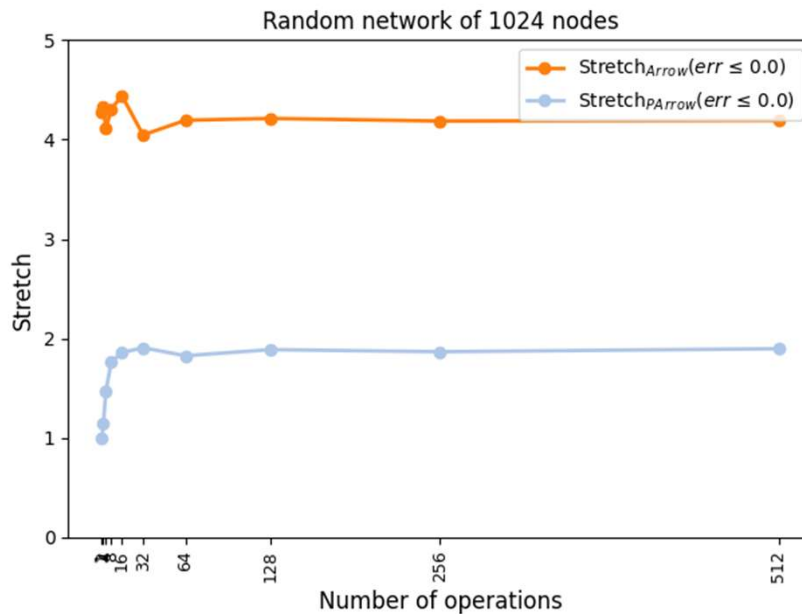
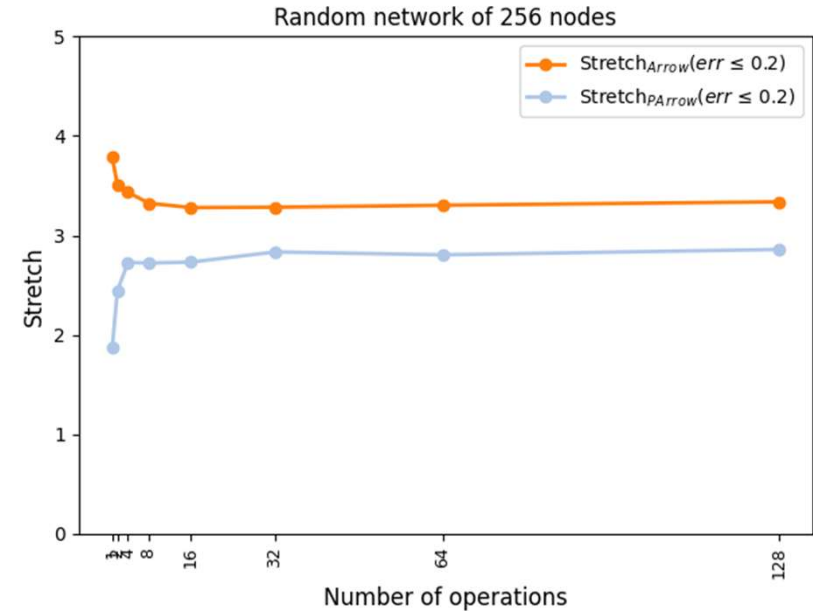
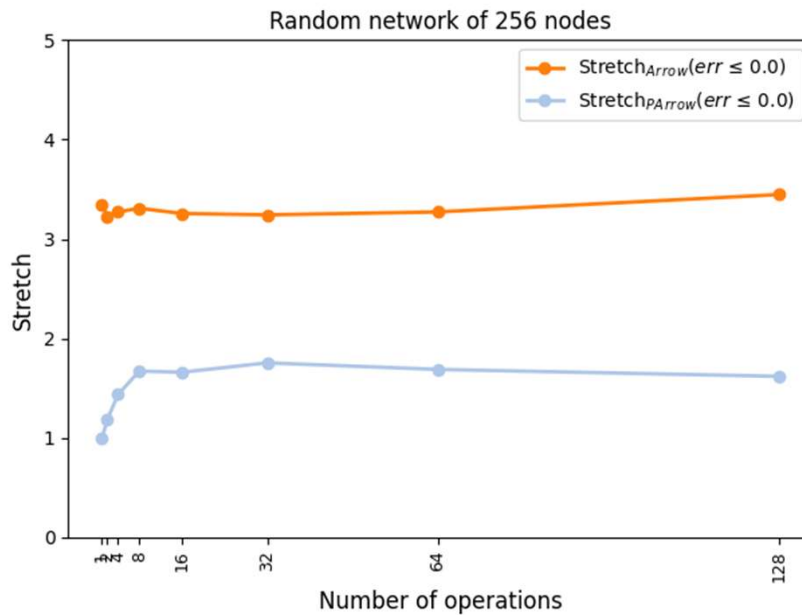
Ideas

- **Tree-based (called Parrow)**
 - Construct Steiner Tree of predicted nodes
 - Use existing ideas, many protocols
 - Augment Steiner Tree with nodes not already in the tree
 - Consider Steiner Tree as a super node
 - Construct a MST of super node and nodes not in Steiner Tree
- Prove the claimed bounds

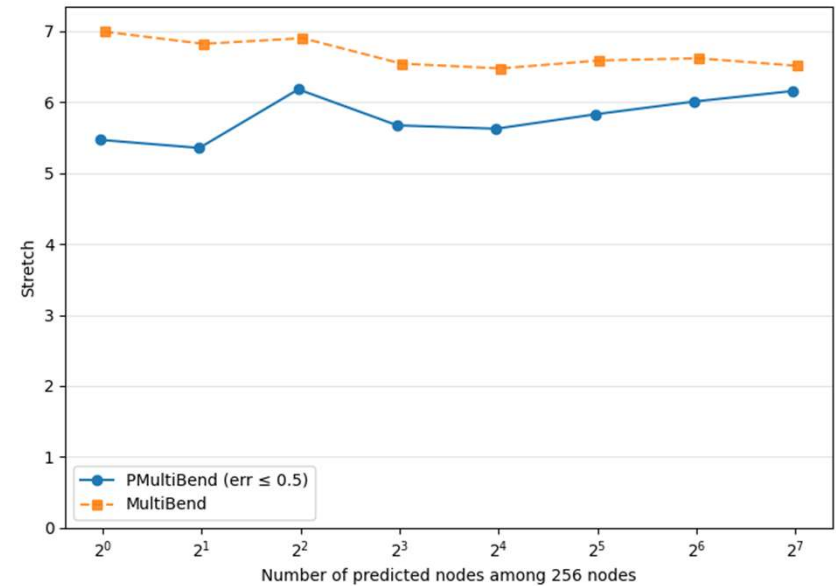
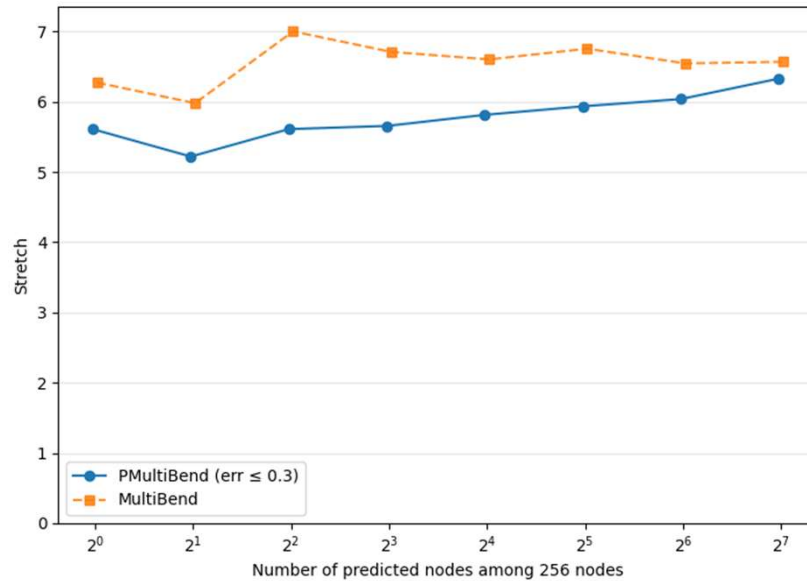
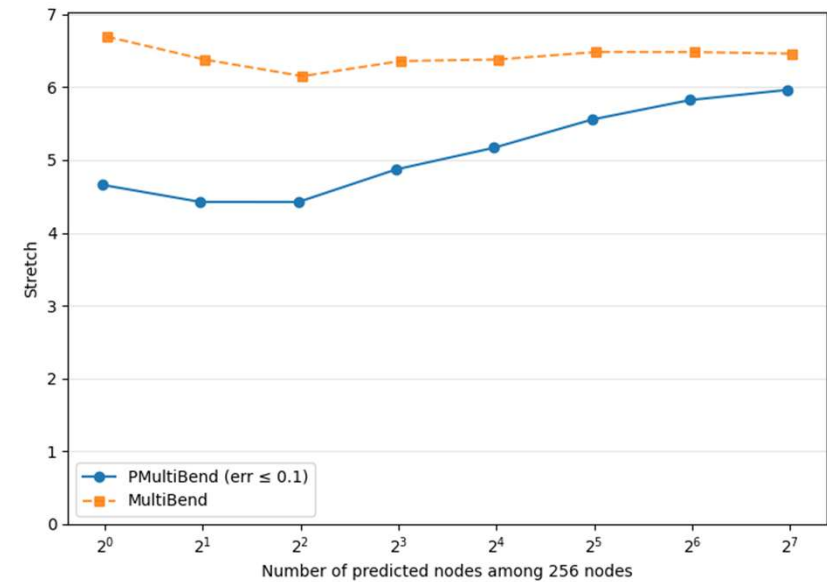
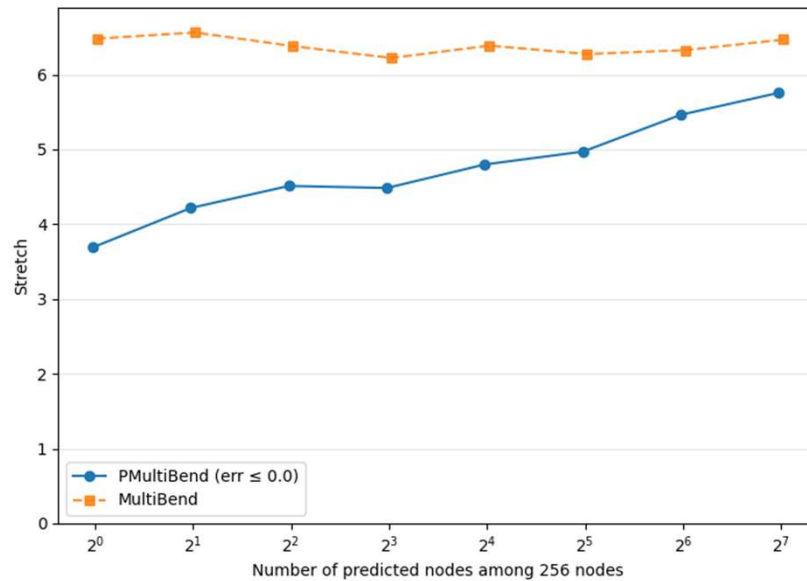
Ideas

- **Cluster-based (called PSpiral)**
 - Use the cluster hierarchy of existing protocols
 - Show that no major change is required,
 - except the object needs to be published up to the root cluster to account for predictor error
 - Prove the claimed bounds

Experimental Results (Tree-Based)



Experimental Results (Cluster-Based)



Summary

Two protocols for accessing a shared object in the network in the location prediction model:

- Tree-based
- Cluster-based

Future Directions

- Online Location Prediction Model - predicted locations released over time
- Other prediction models (Time?)
- Detailed Experiments

Thank you!!!